

Anonymity on Quicksand

Using BGP to compromise Tor



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HotNets

October, 28 2014

Joint work with

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**WE ARE
ANONYMOUS**



**WE ARE NOT
ANONYMOUS FOR LONG.**



**WE ARE NOT
ANONYMOUS FOR LONG.
COURTESY OF BGP.**

Internet communications are *not* anonymous

Looking at an Internet communication, one can

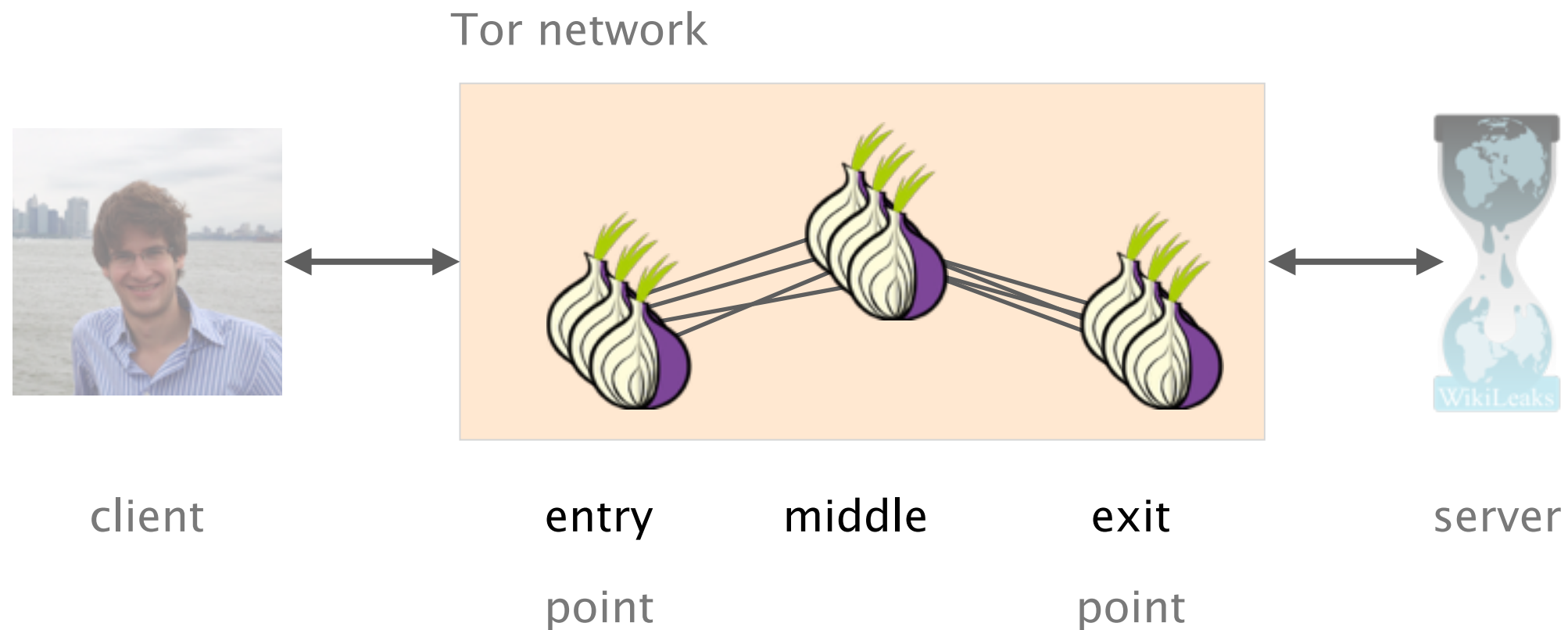
- infer who is talking to whom
- infer physical locations
- use that to track behavior and interests

even if the communication is encrypted

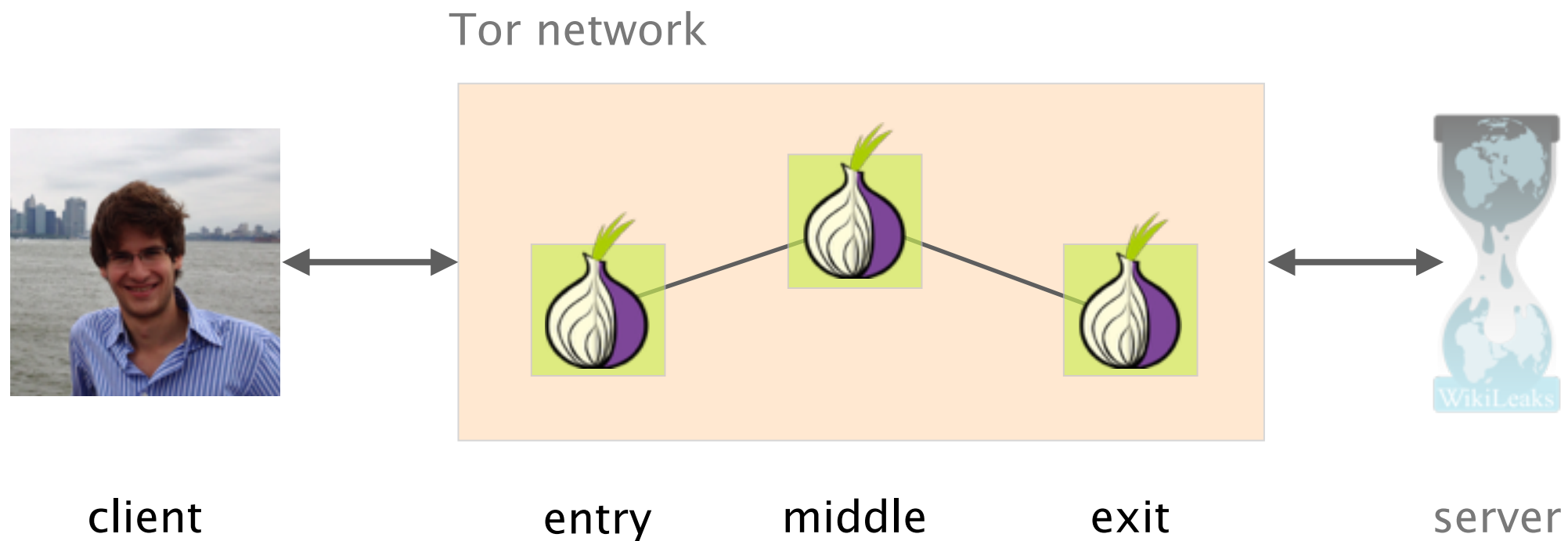
Tor aims at preventing adversaries to follow packets between a sender and a receiver



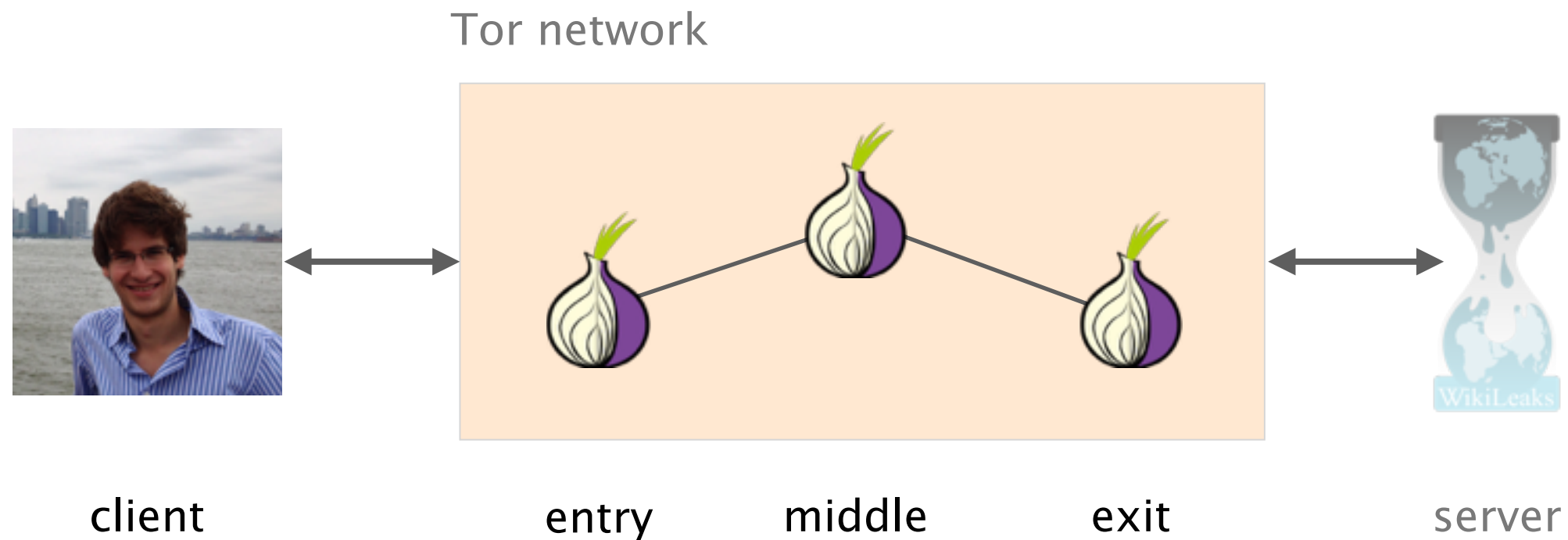
To do that,
Tor bounces traffic around a network of relays



Tor clients start by selecting
3 relays, one of each type



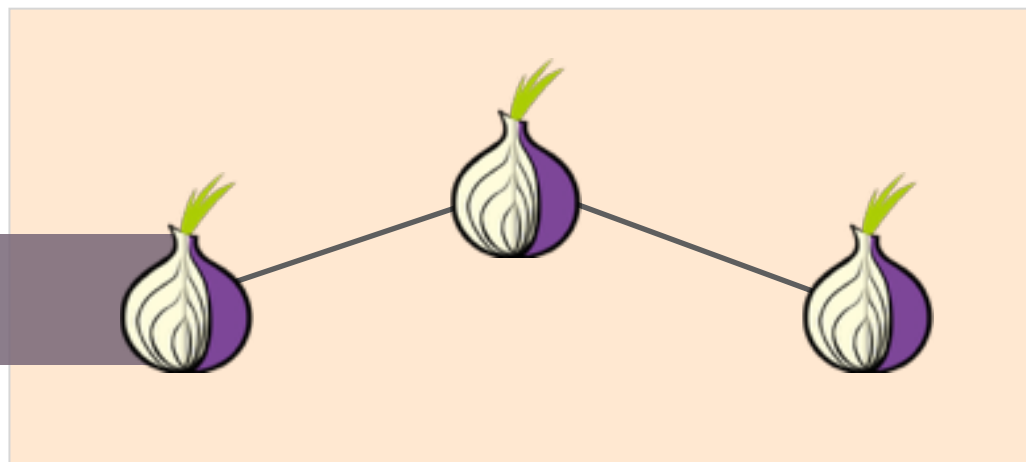
Tor clients then incrementally
build encrypted circuits through them



Tor network



client



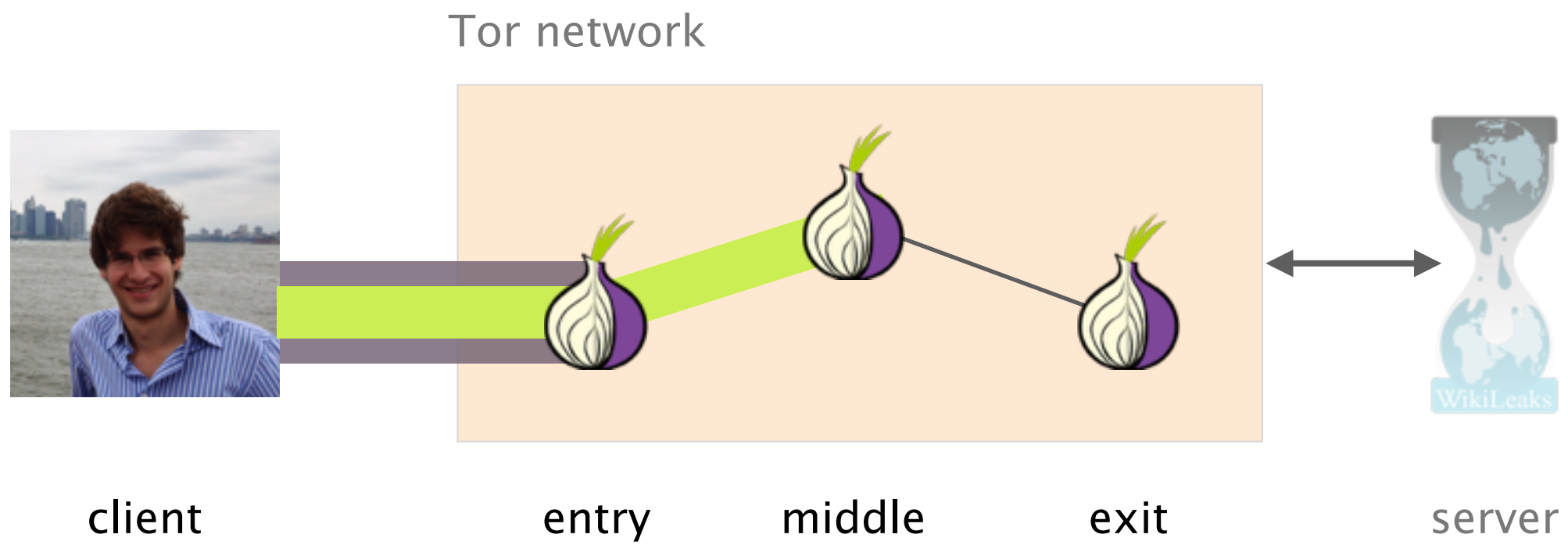
entry

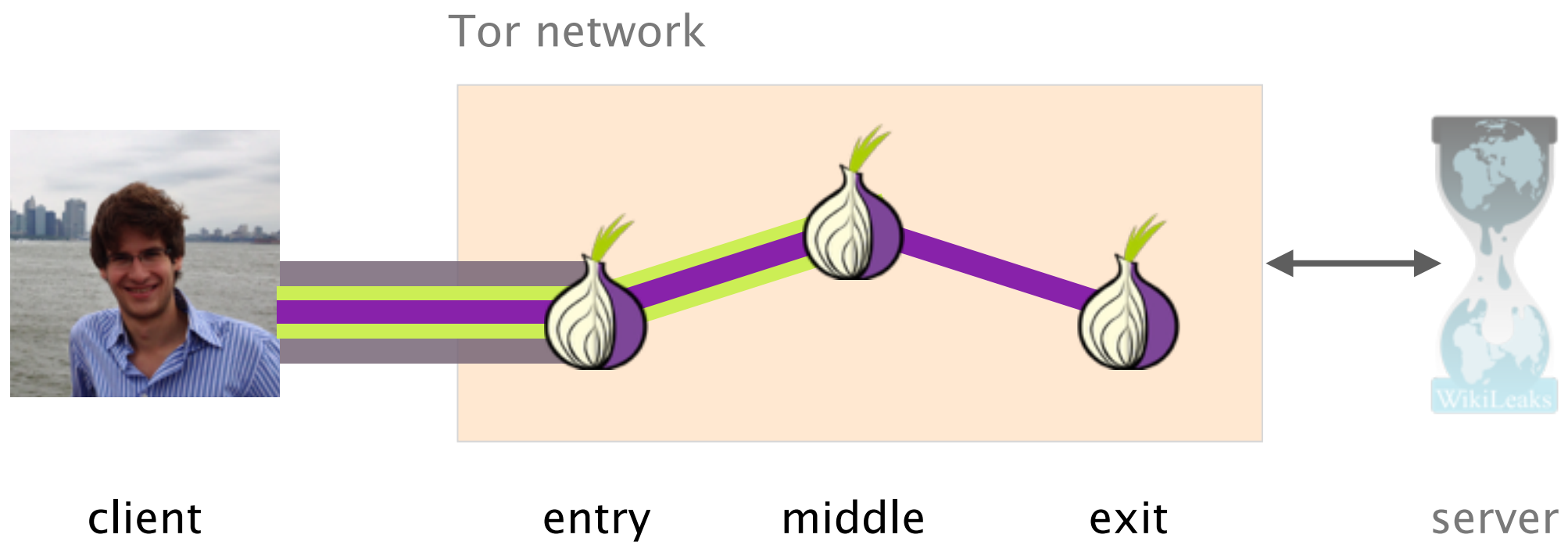
middle

exit

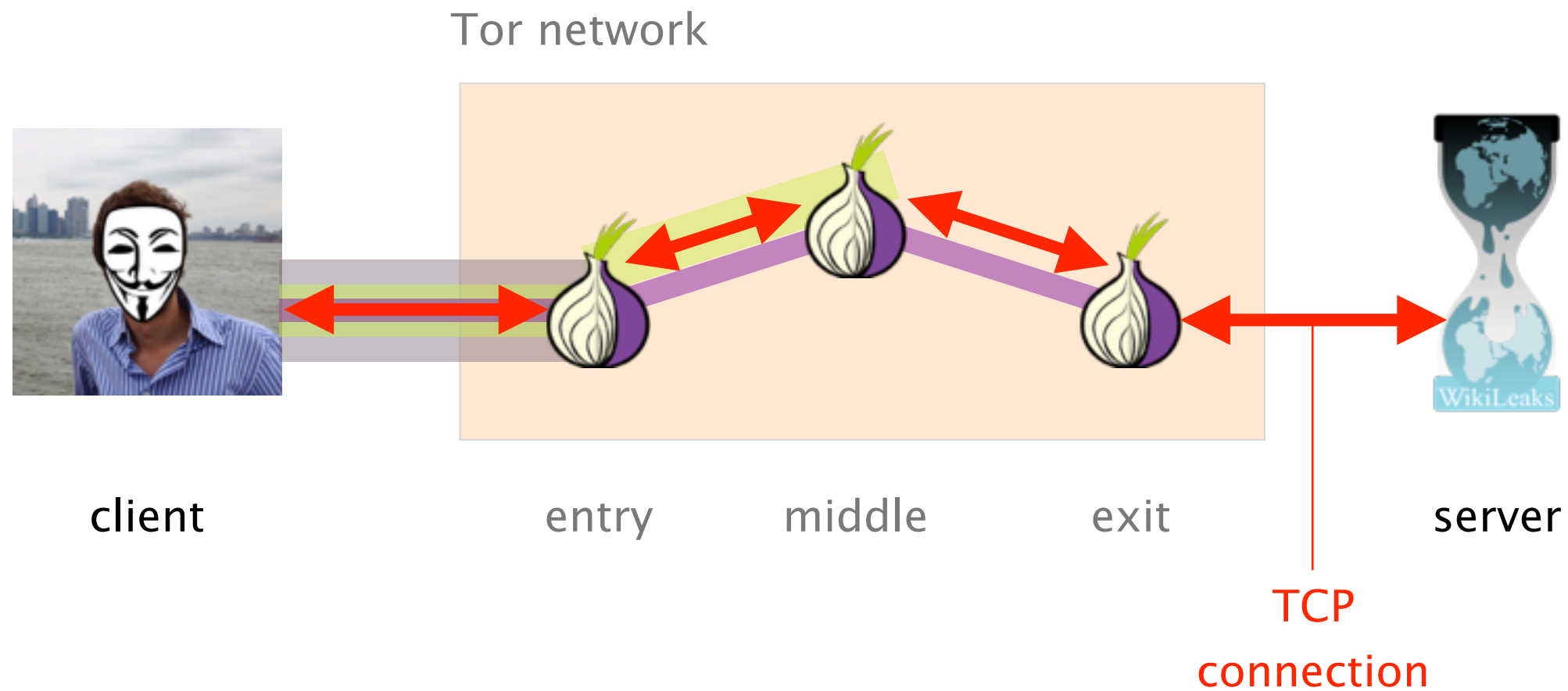


server

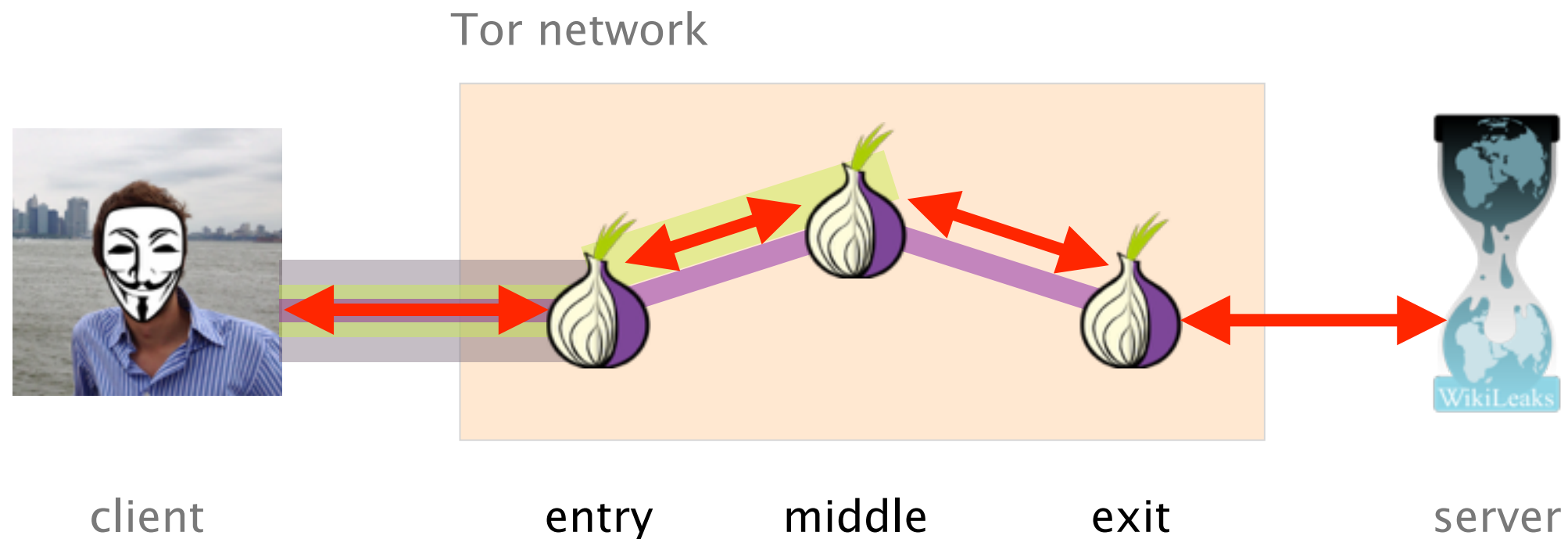


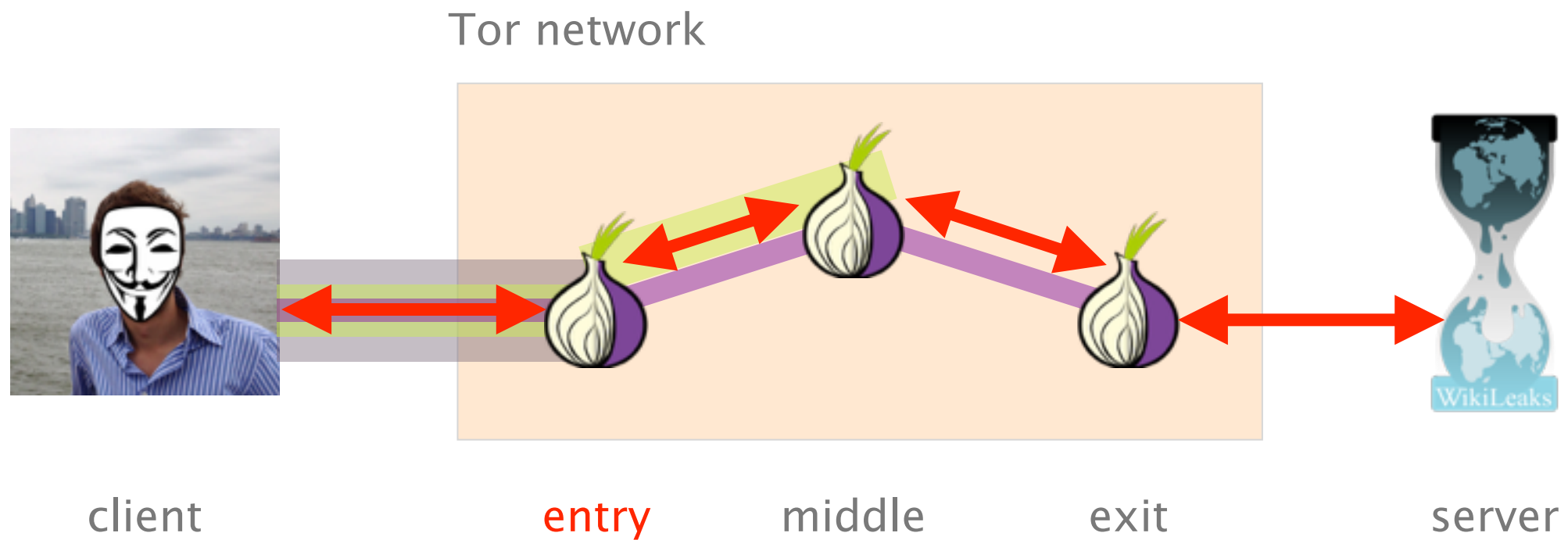


Anonymous communication takes place
by forwarding across consecutive tunnels

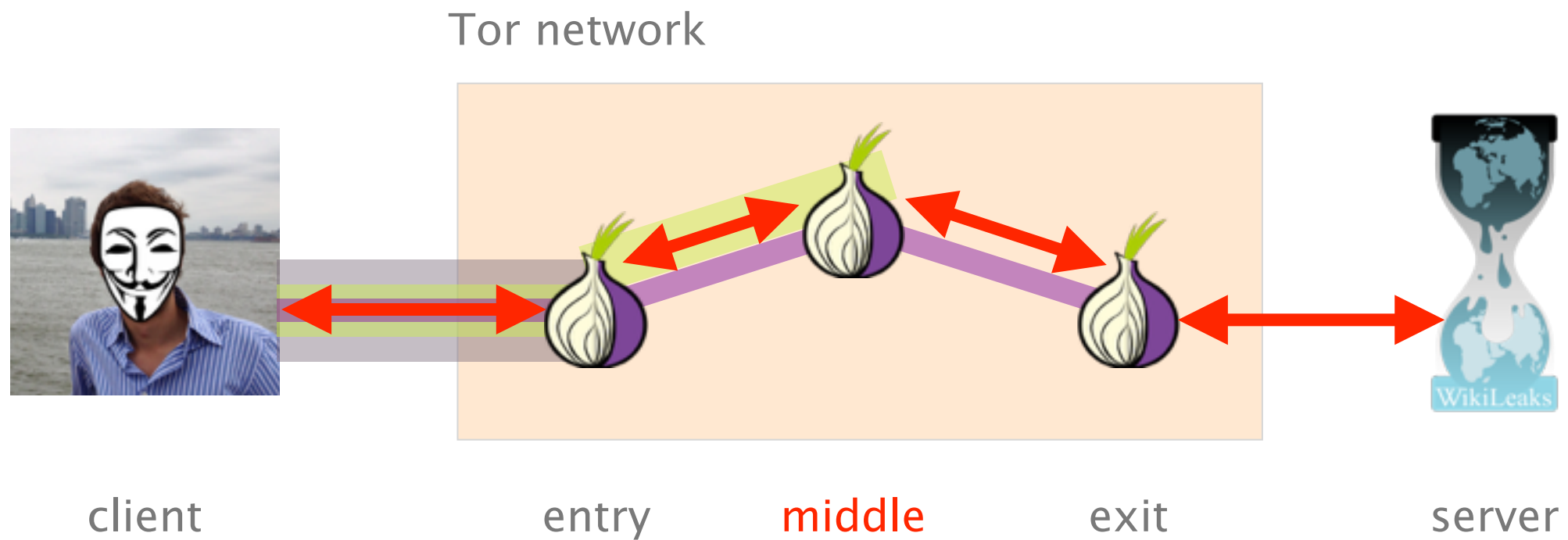


Not a single Tor entity knows
the association (client, server)

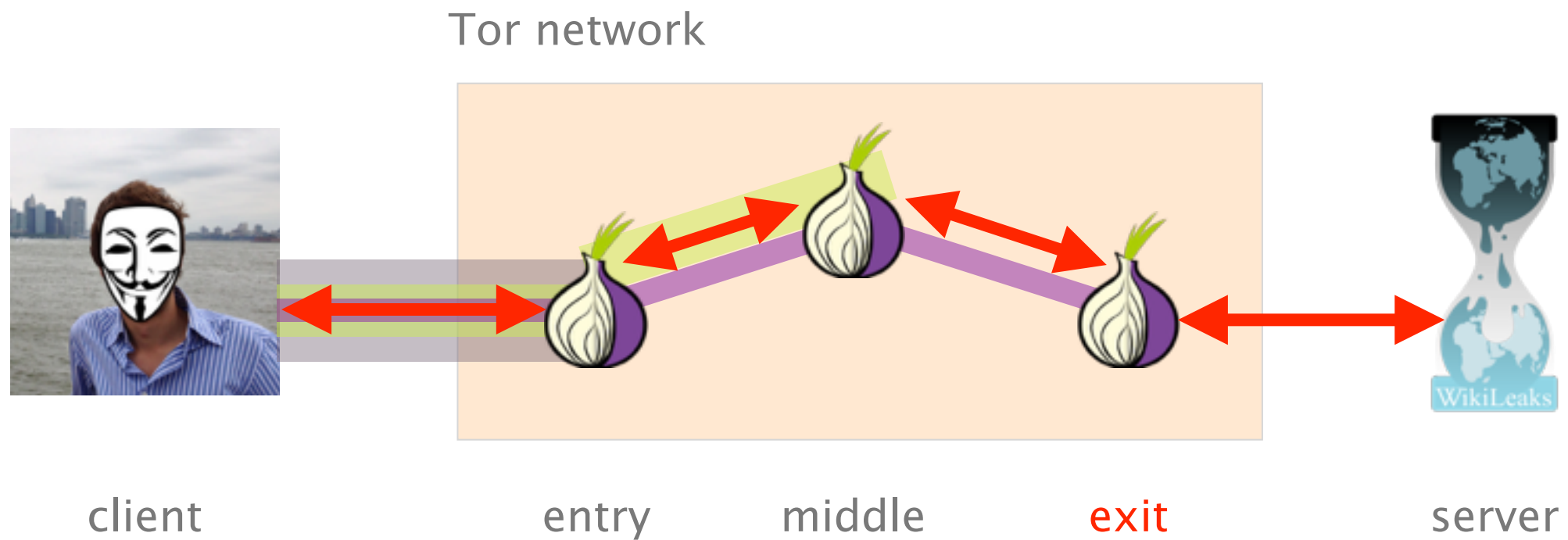




knows the source,
not the destination

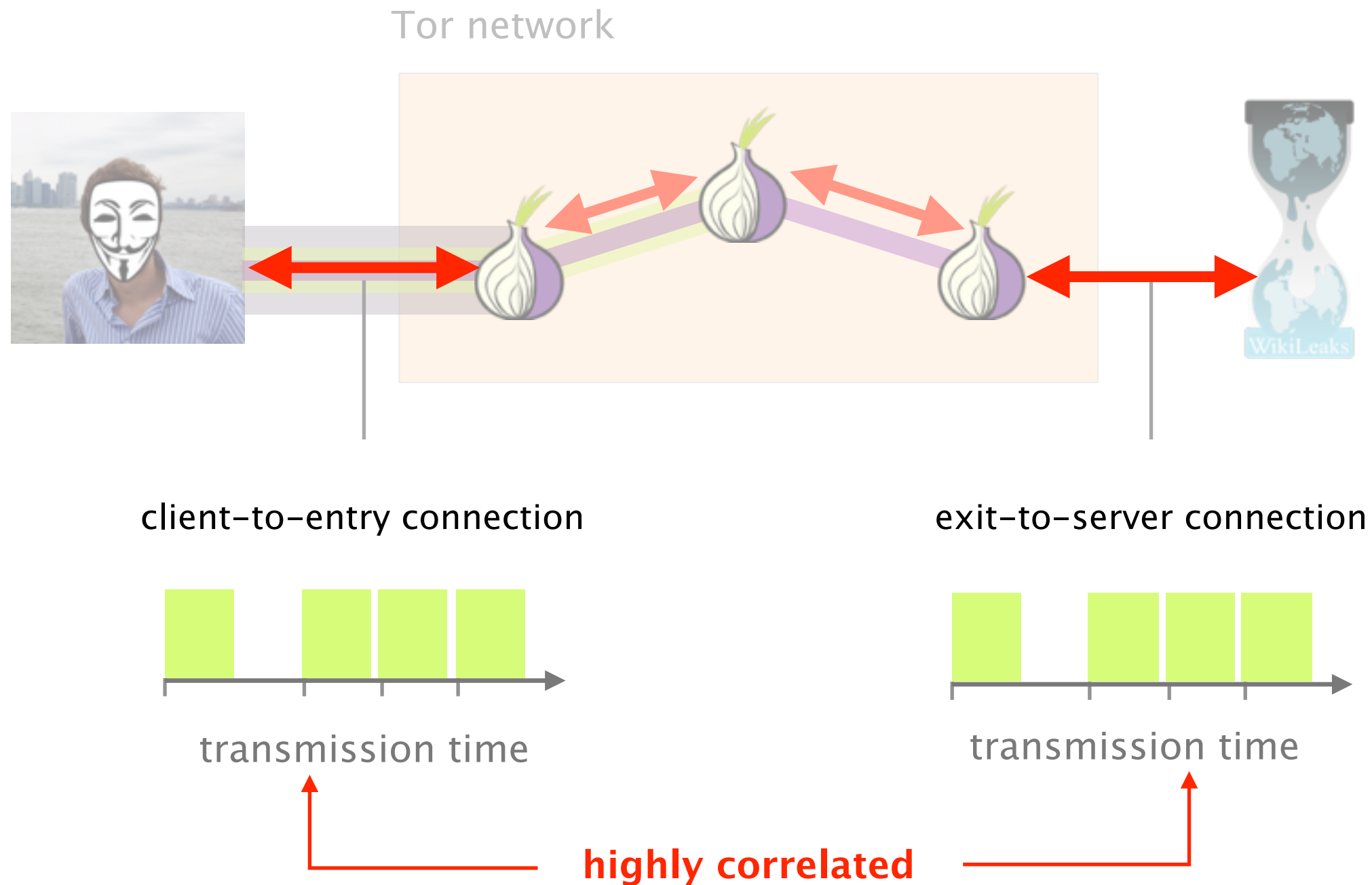


knows neither the source,
nor the destination



knows the destination,
nor the source

Traffic entering and leaving Tor is highly correlated



By correlating client-to-entry & exit-to-server flows,
one can effectively de-anonymize Tor users

Traffic correlation attacks require to see
client-to-entry and exit-to-server traffic

Traffic correlation attacks require to **see**
client-to-entry and exit-to-server traffic



How?

Two ways

Manipulate Tor
malicious relay

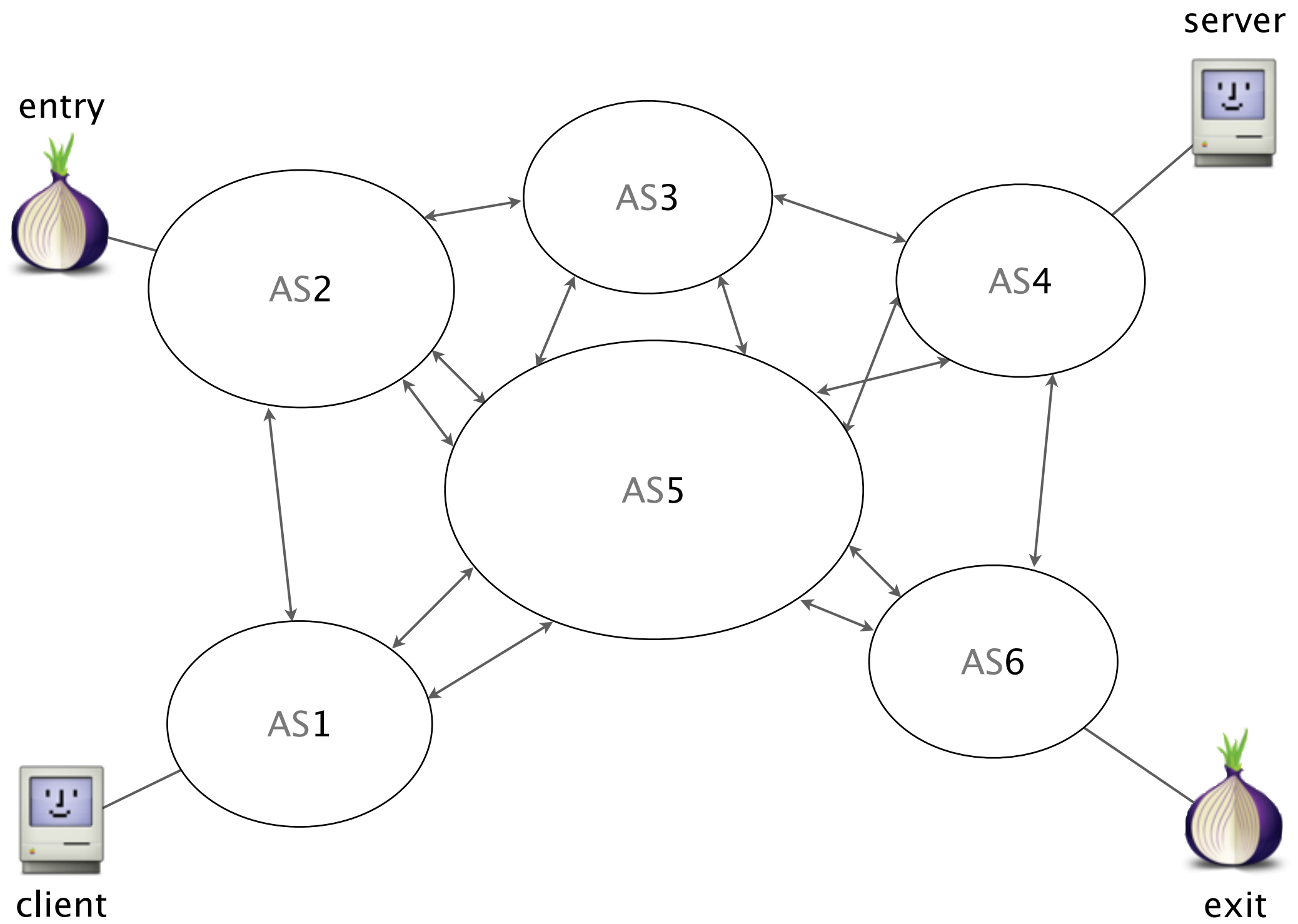
Manipulate routing
malicious networks

Two ways

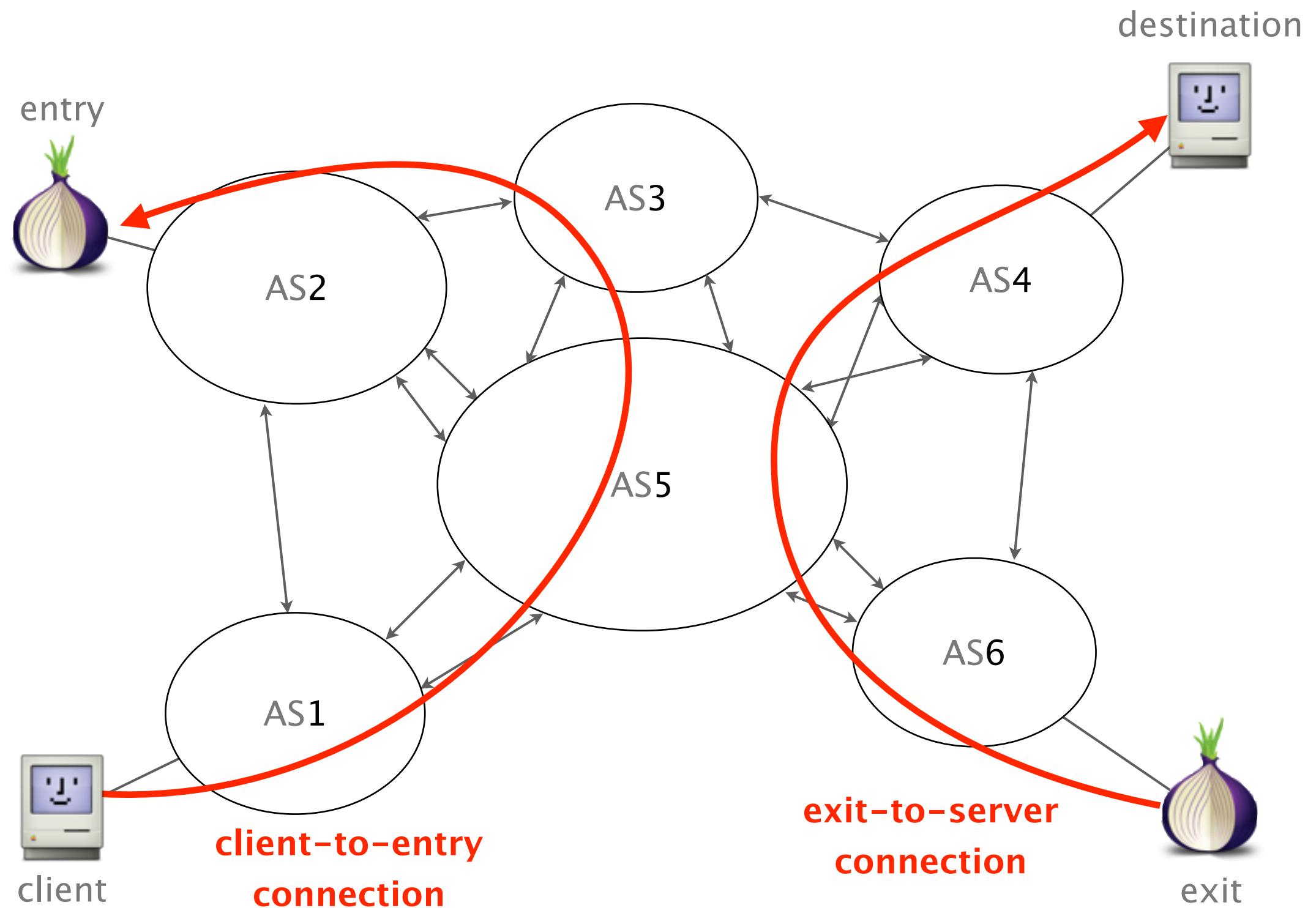
Manipulate Tor
malicious relay

Manipulate routing
malicious networks

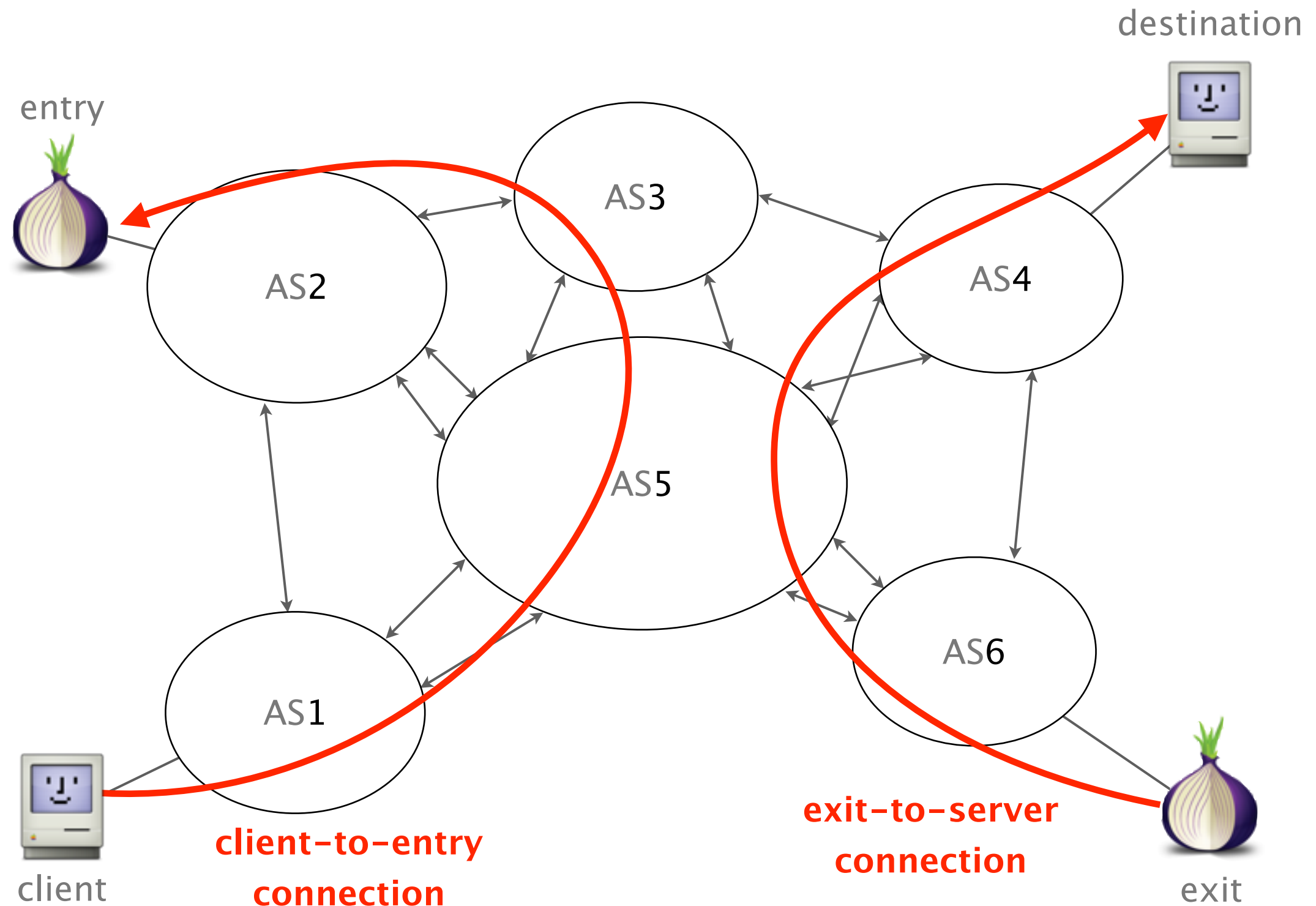
This talk

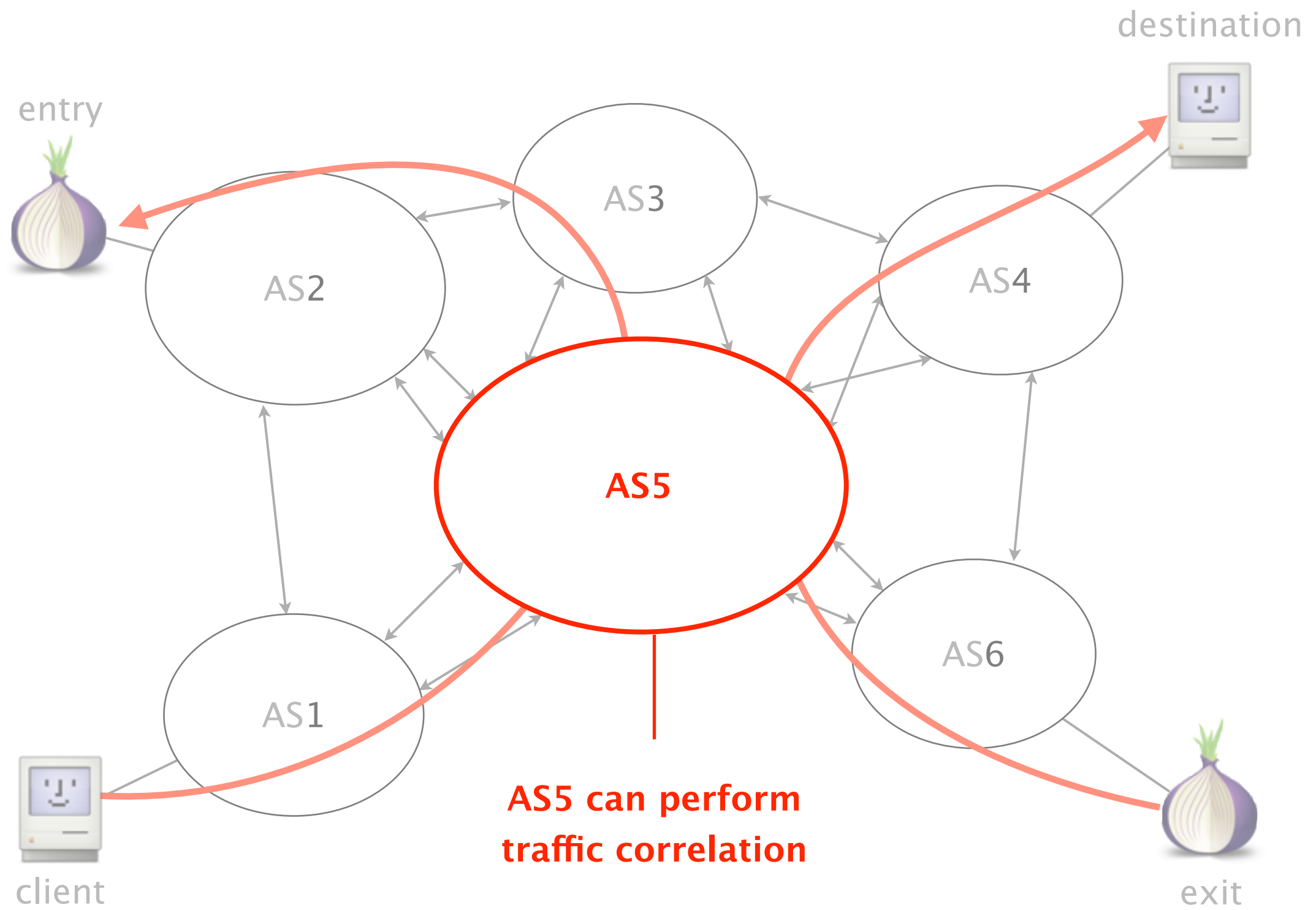


Tor connections get routed according to BGP



Traffic correlation attacks require to **see**
client-to-entry *and* **exit-to-server** traffic





Network-level adversaries are a known problem

Related work

2004	Location diversity in anonymity networks	Feamster and Dingledine
2007	Sampled traffic analysis by Internet-exchange-level adversaries	Murdoch and Zieliński
2009	AS-awareness in Tor Path Selection	Edman and Syverson
2013	Traffic correlation on Tor by realistic adversaries	Johnson <i>et al.</i>

However, these works assume
that the Internet is **static**

However, these works assume
that the Internet is **static**

... which is **not** the case

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that the Internet is **static**

... which is **not** the case

Contribution

What's the impact on Tor?

User anonymity decreases over time
due to BGP dynamics

User anonymity decreases over time due to BGP dynamics

3 BGP-induced
causes

Natural BGP convergence
policy changes, failures, etc.

Active BGP manipulation
IP prefix hijack, interception (MITM), etc.

Asymmetric routing
seeing one direction of the connection is enough

Anonymity on Quicksand

Using BGP to compromise Tor



- 1 **Attacks**
All your traffic belongs to me
- 2 **Preliminary results**
Eyes wide open
- 3 **Countermeasures**
Close the curtains

Anonymity on Quicksand

Using BGP to compromise Tor



1

Attacks

All your traffic belongs to me

Preliminary results

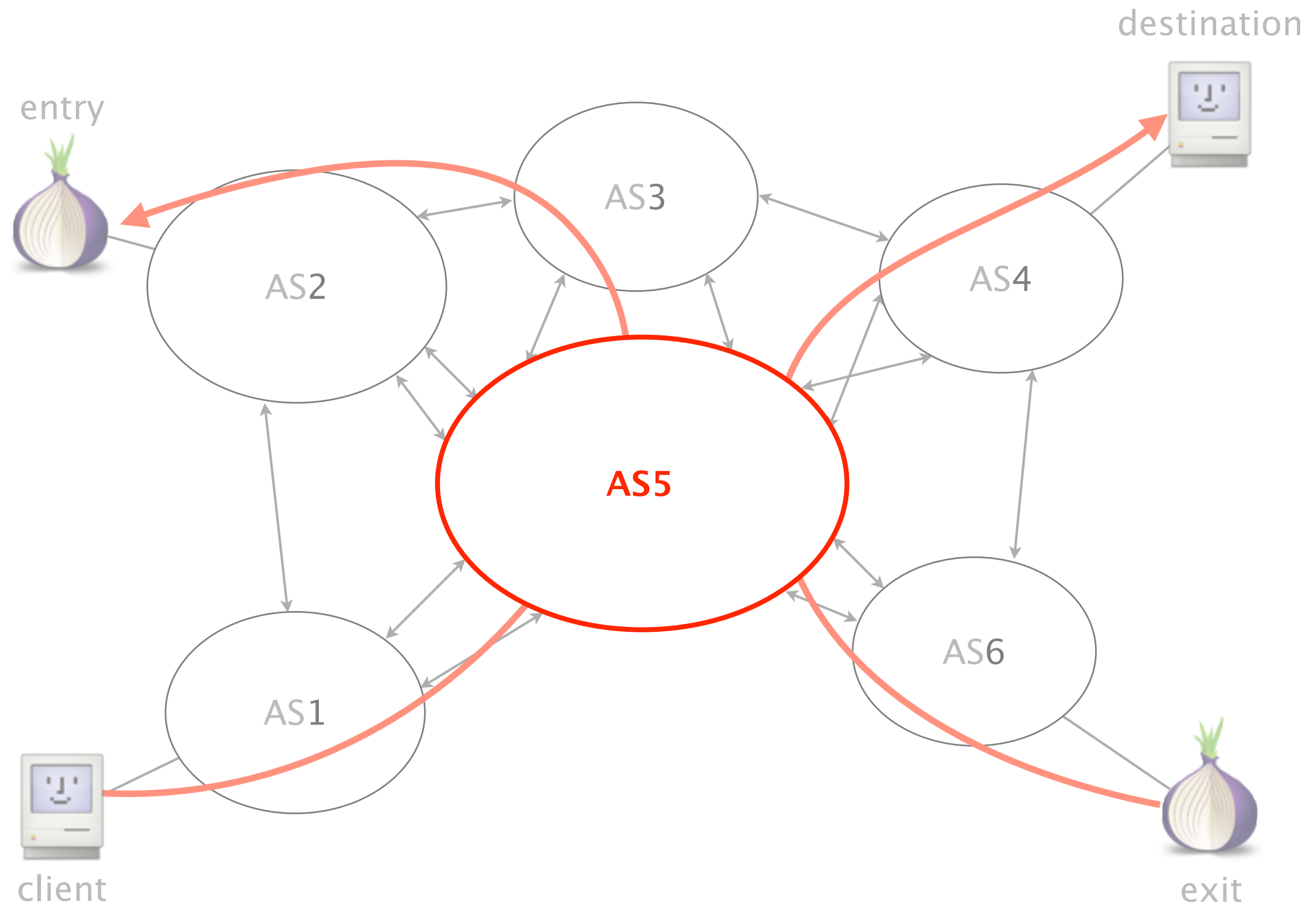
Eyes wide open

Countermeasures

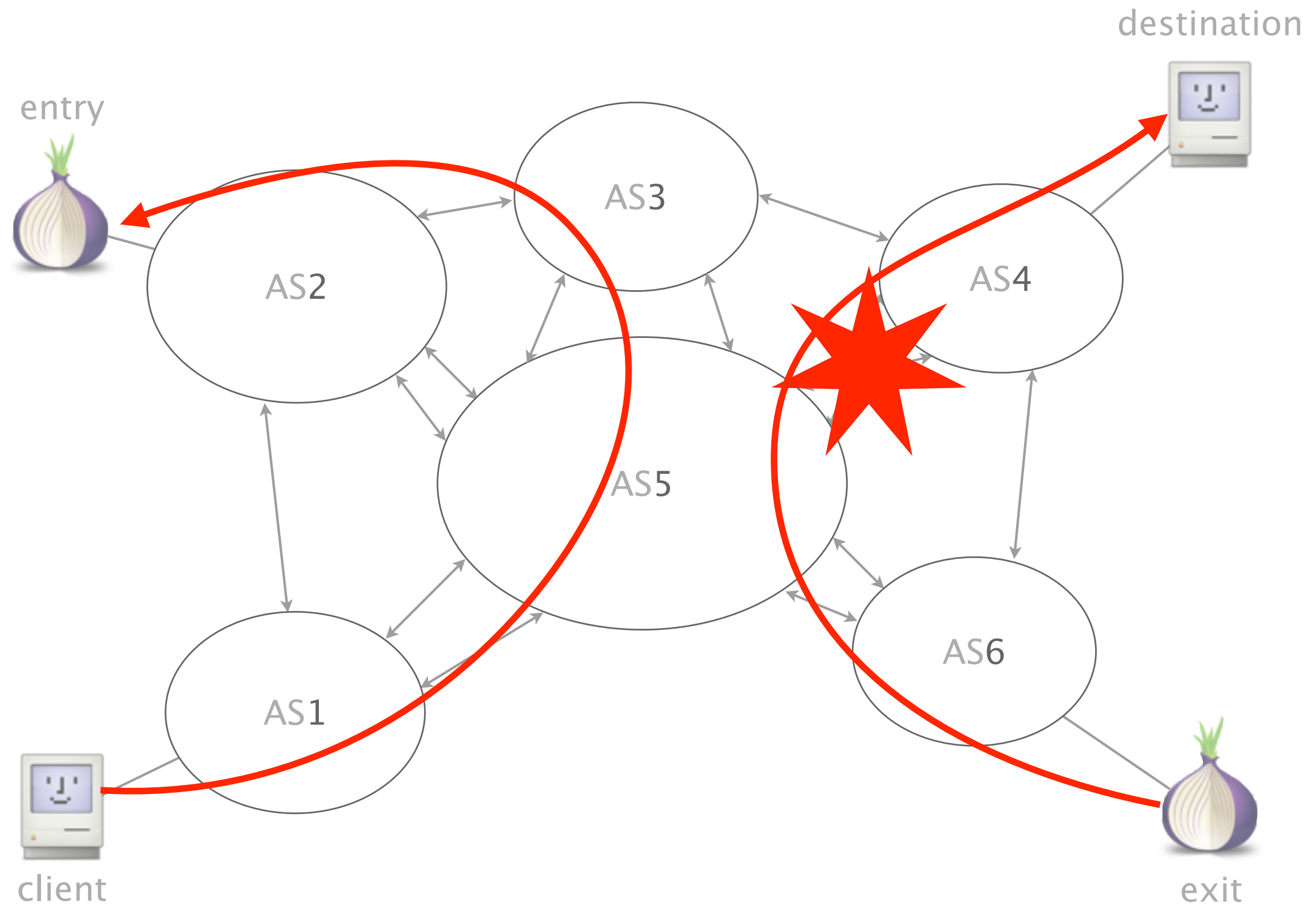
Close the curtains

Attack#1: Natural BGP dynamics increases
the number of AS-level adversaries

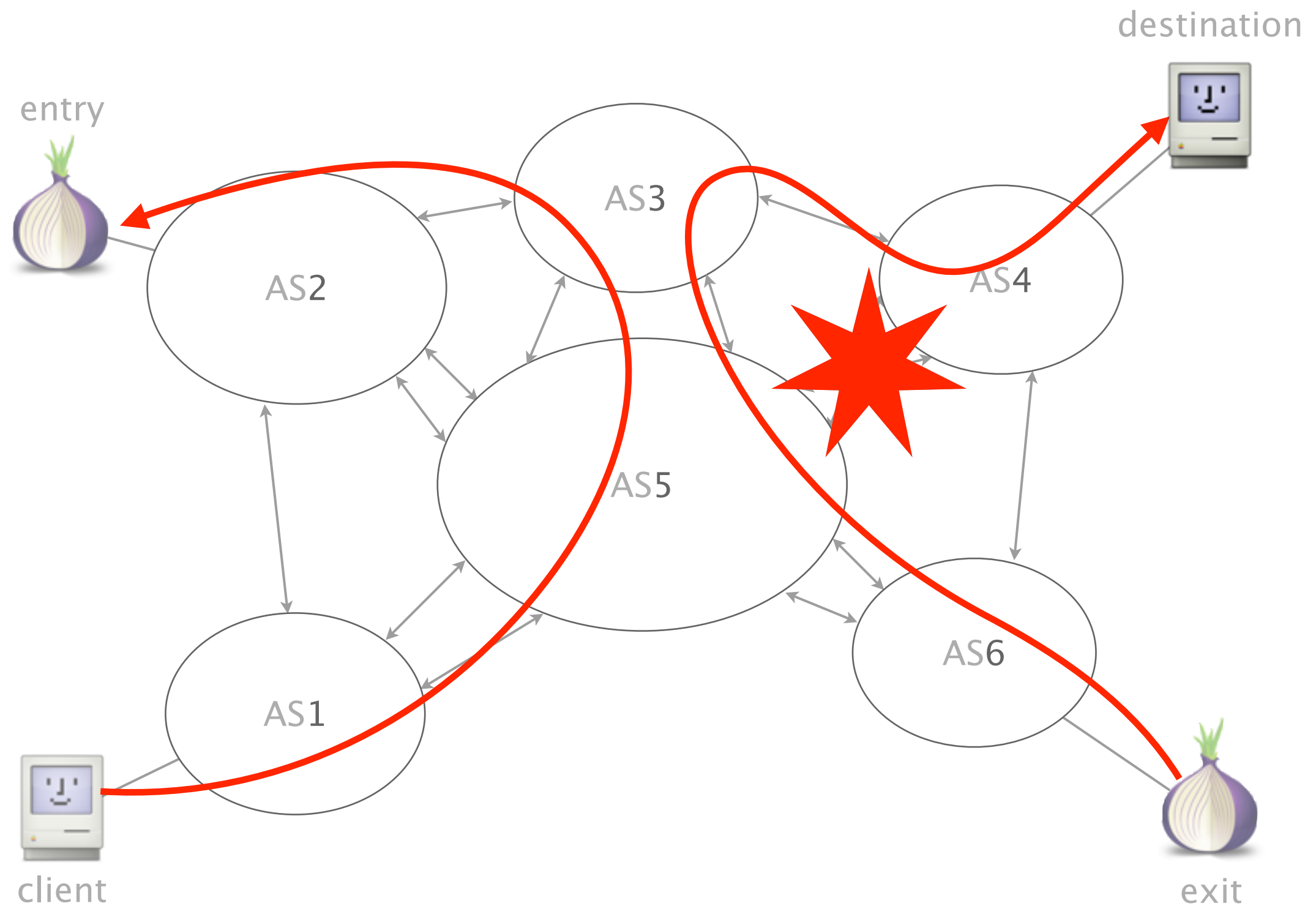
Initially, only AS5 is seeing traffic
client-to-entry and exit-to-server traffic



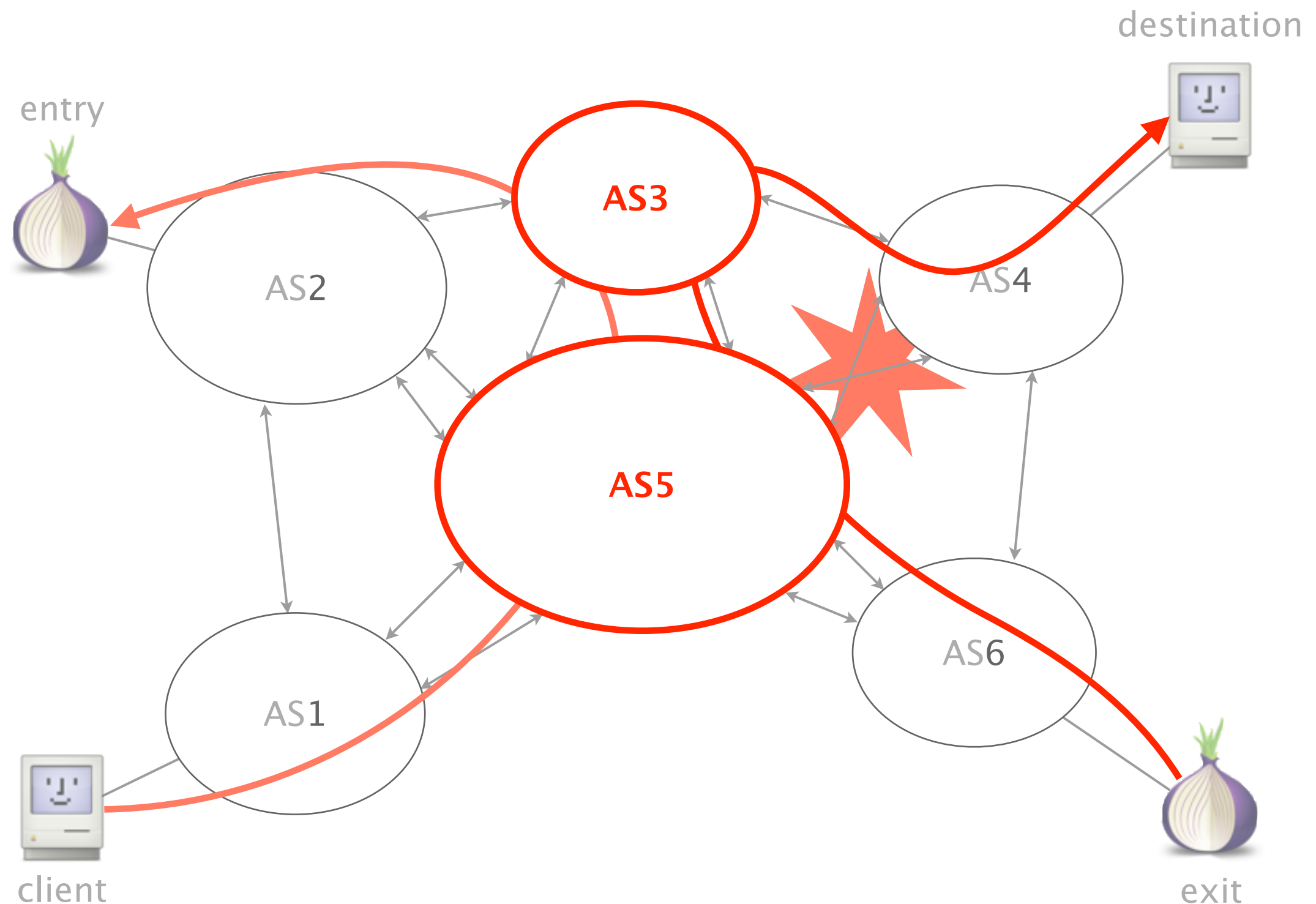
Link between AS4 and AS5 fails



Traffic gets rerouted via AS3

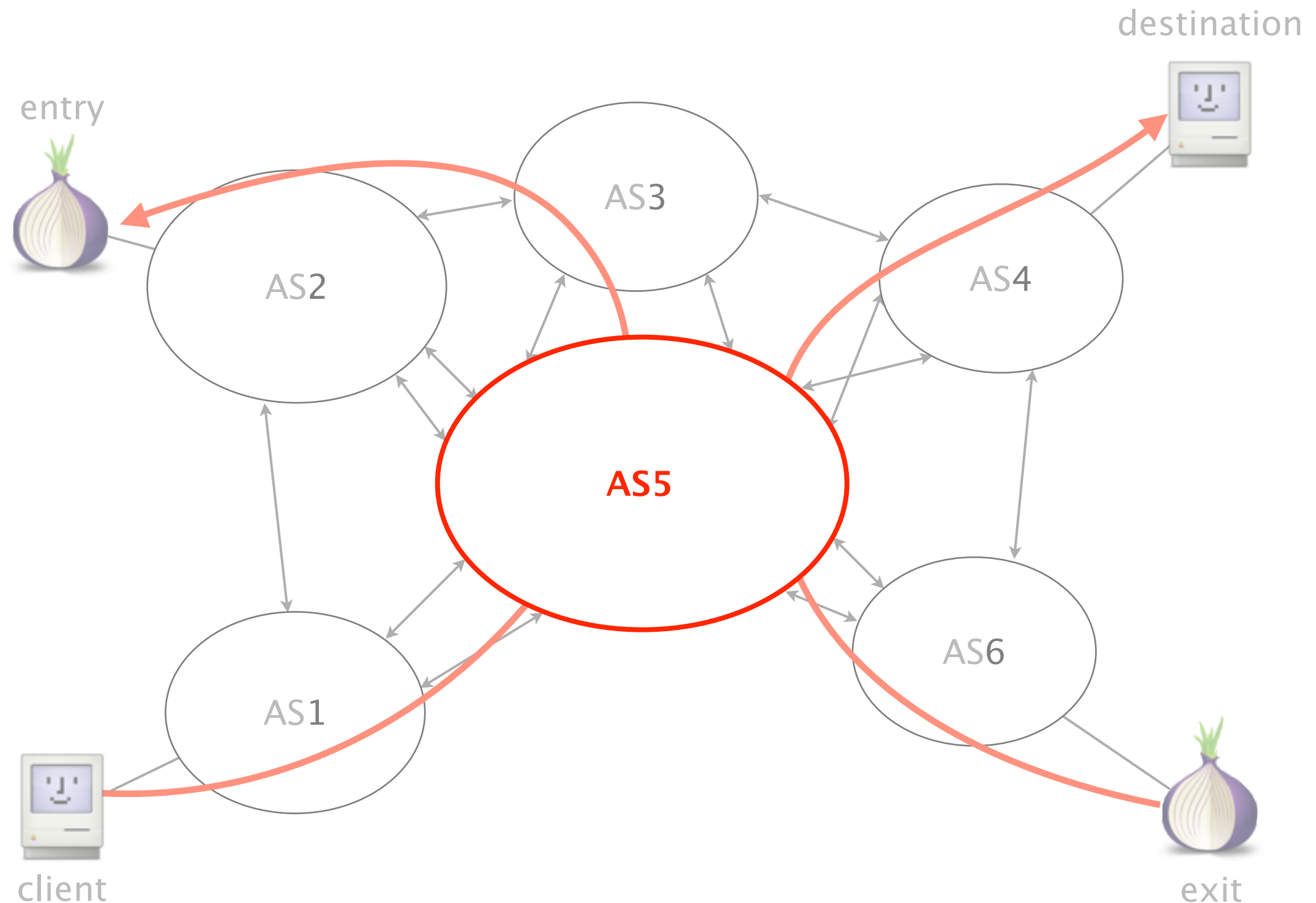


Now, both AS3 and AS5 are seeing
client-to-entry and exit-to-server traffic

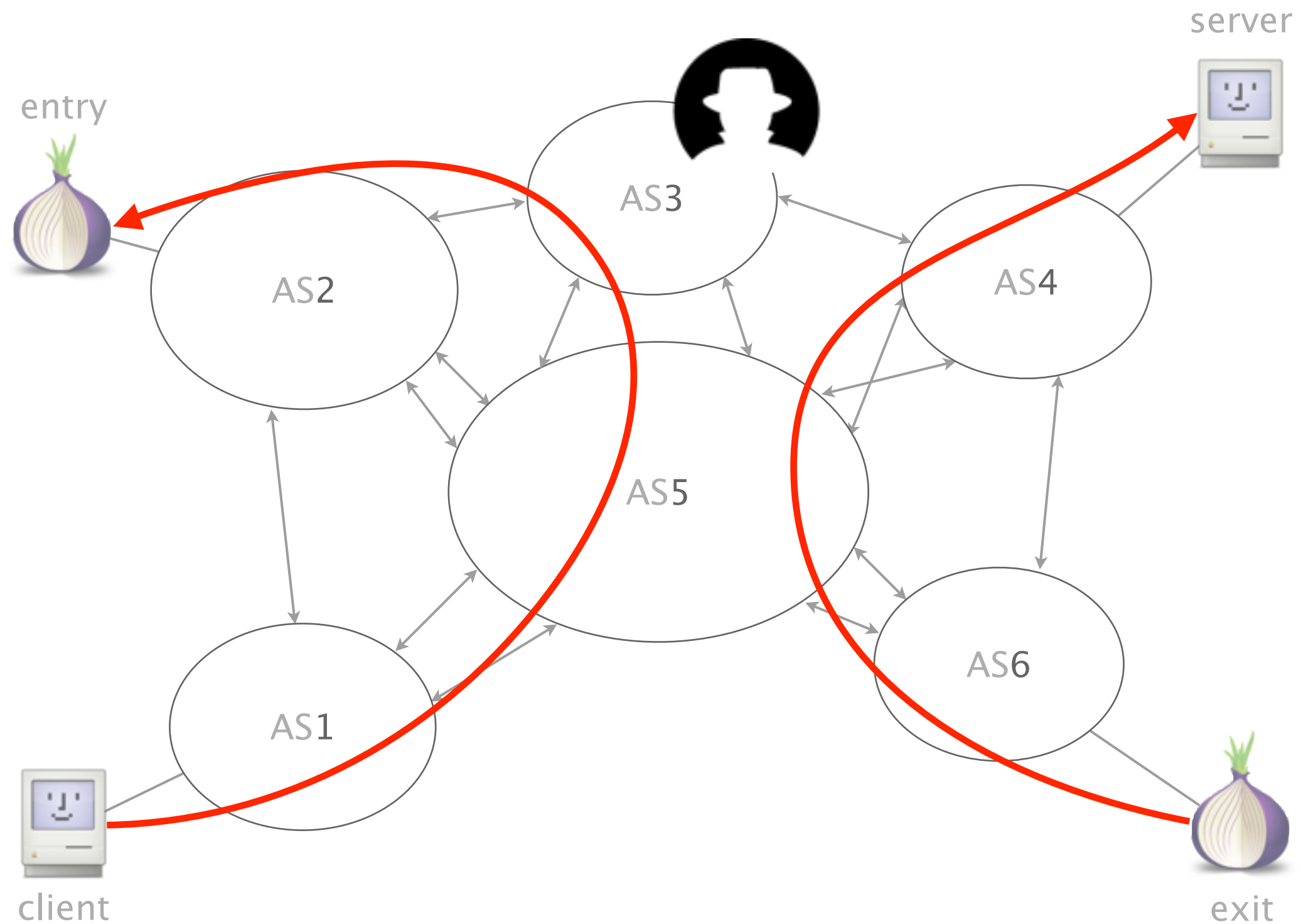


Attack#2: BGP hijacking attacks enable on-demand, fine-grained Tor attacks

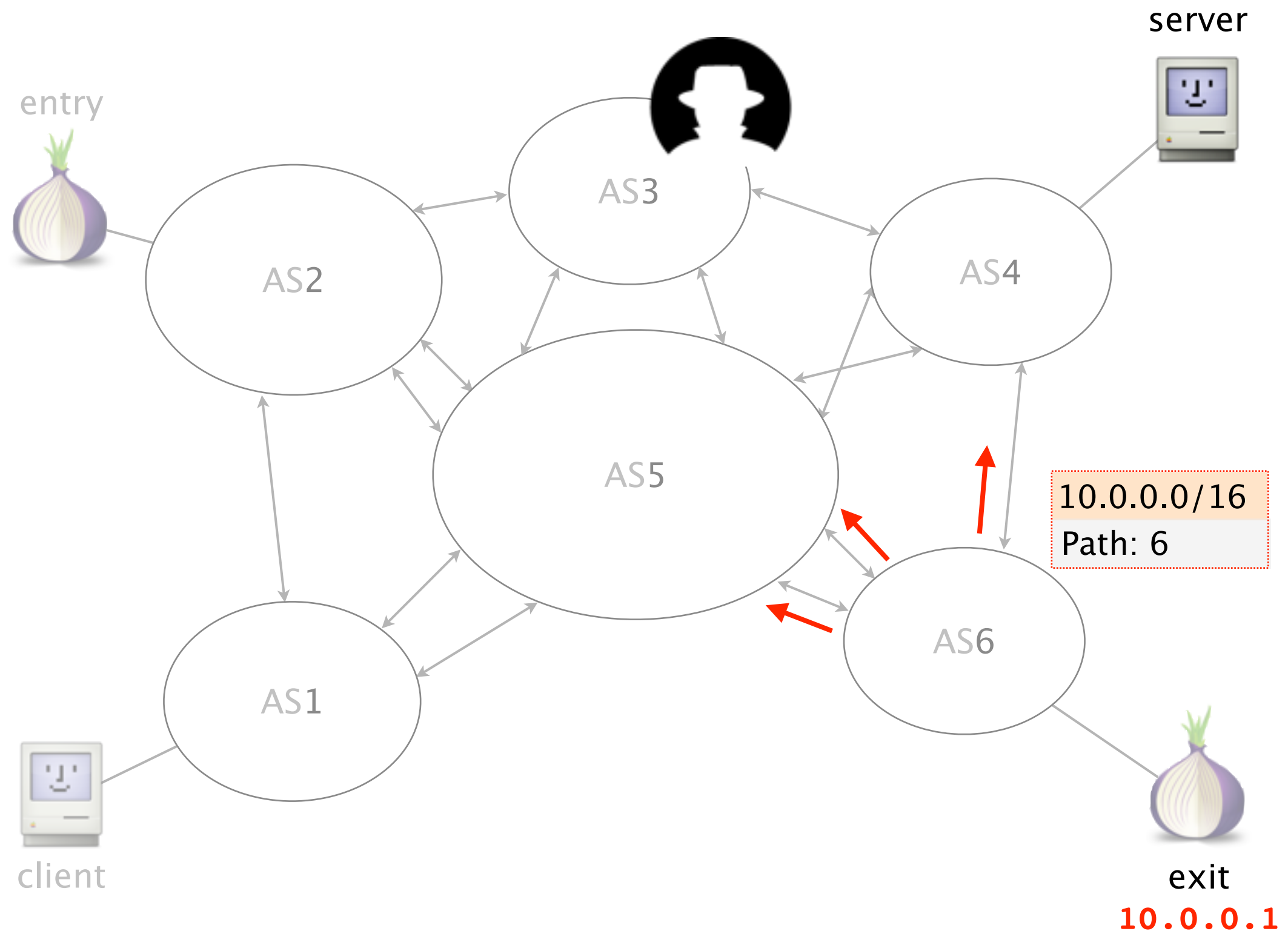
Initially, only AS5 is seeing traffic entering and exiting the Tor network



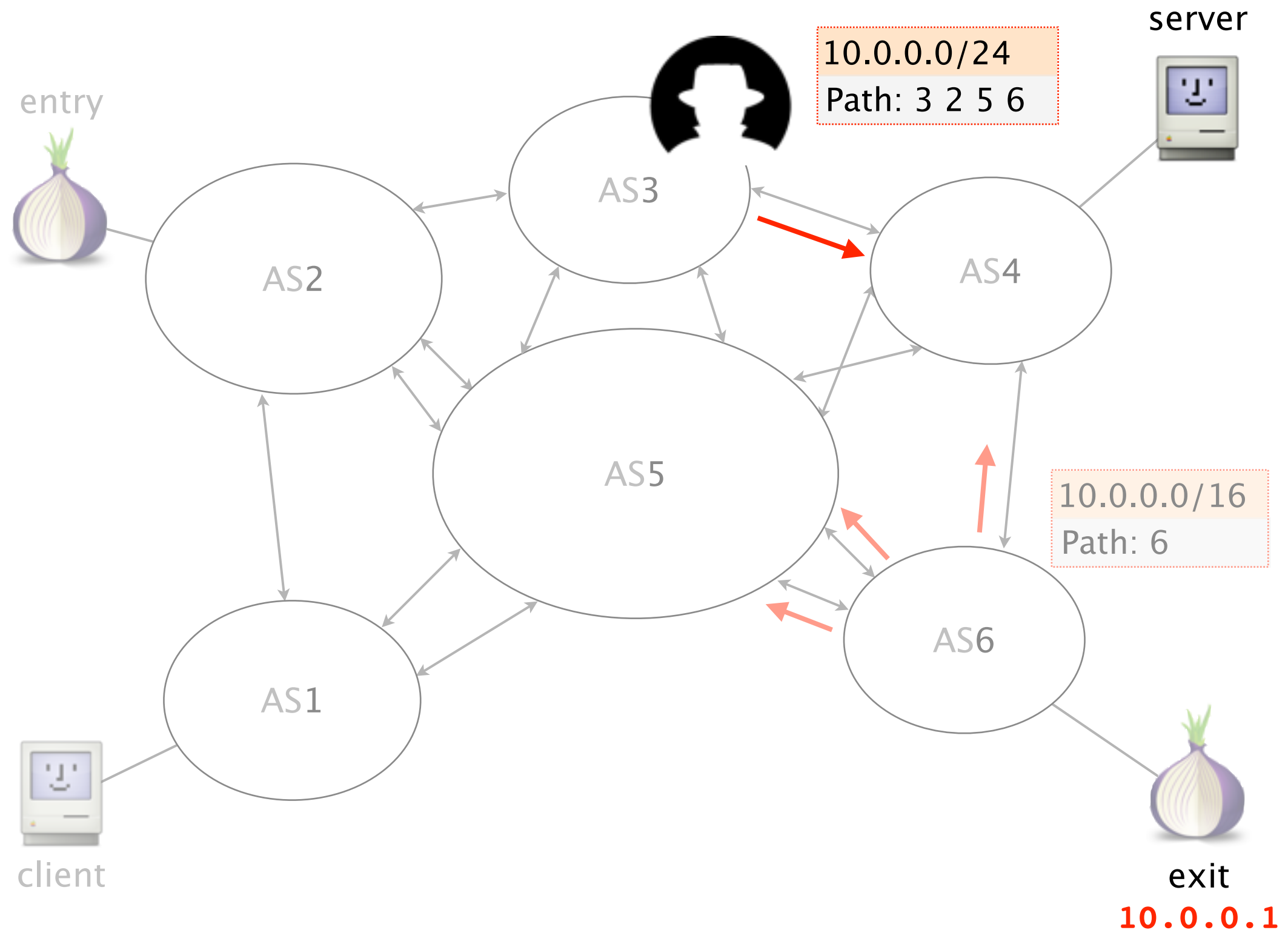
Assume that AS3 is a malicious AS,
and wants to observe Tor traffic

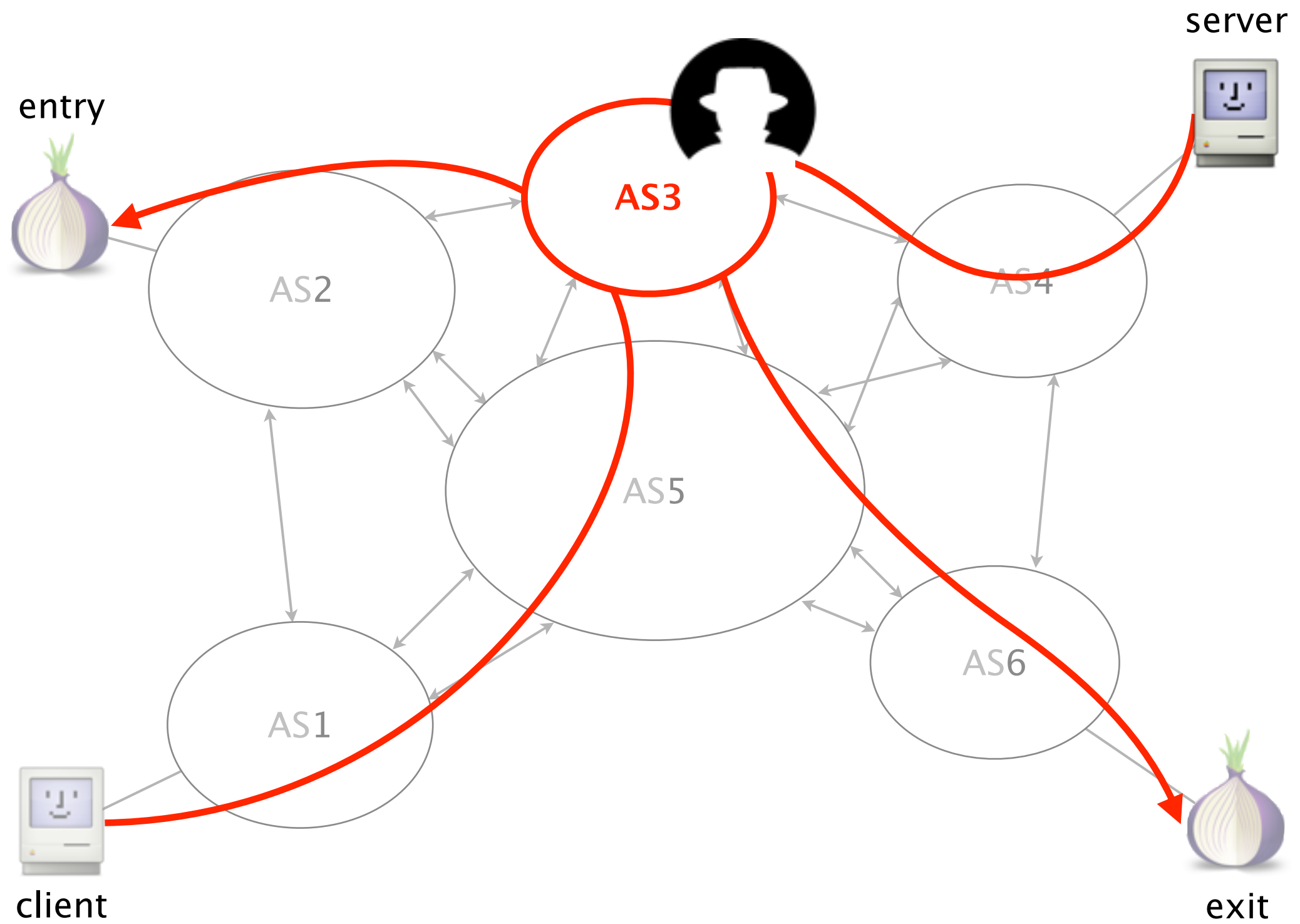


AS3 can put itself on server-to-exit paths by hijacking Tor prefixes



AS3 can put itself on server-to-exit paths by hijacking Tor prefixes





In November 2010,
China Telecom hijacked 50k prefixes during ~20 min

NOVEMBER 18, 2010 COMMENTS (10) VIEWS: 778 ENGINEERING, INTERNET, POLITICS, SECURITY JIM COWIE

China's 18-Minute Mystery

Twitter Facebook Google+ LinkedIn YouTube



When the US-China Economic and Security Review Commission released its report to Congress this week, something slightly unusual happened: *people read it*. And there, buried on [pages 236-247](#), a mystery was revealed, and the media have greedily amplified it.

Did China's government really divert 15% of the Internet's traffic for eighteen minutes in April, effortlessly intercepting sensitive traffic in flight, and generally creating a massively embarrassing [man-in-the-middle attack](#) on vulnerable global communications?

China Telecom

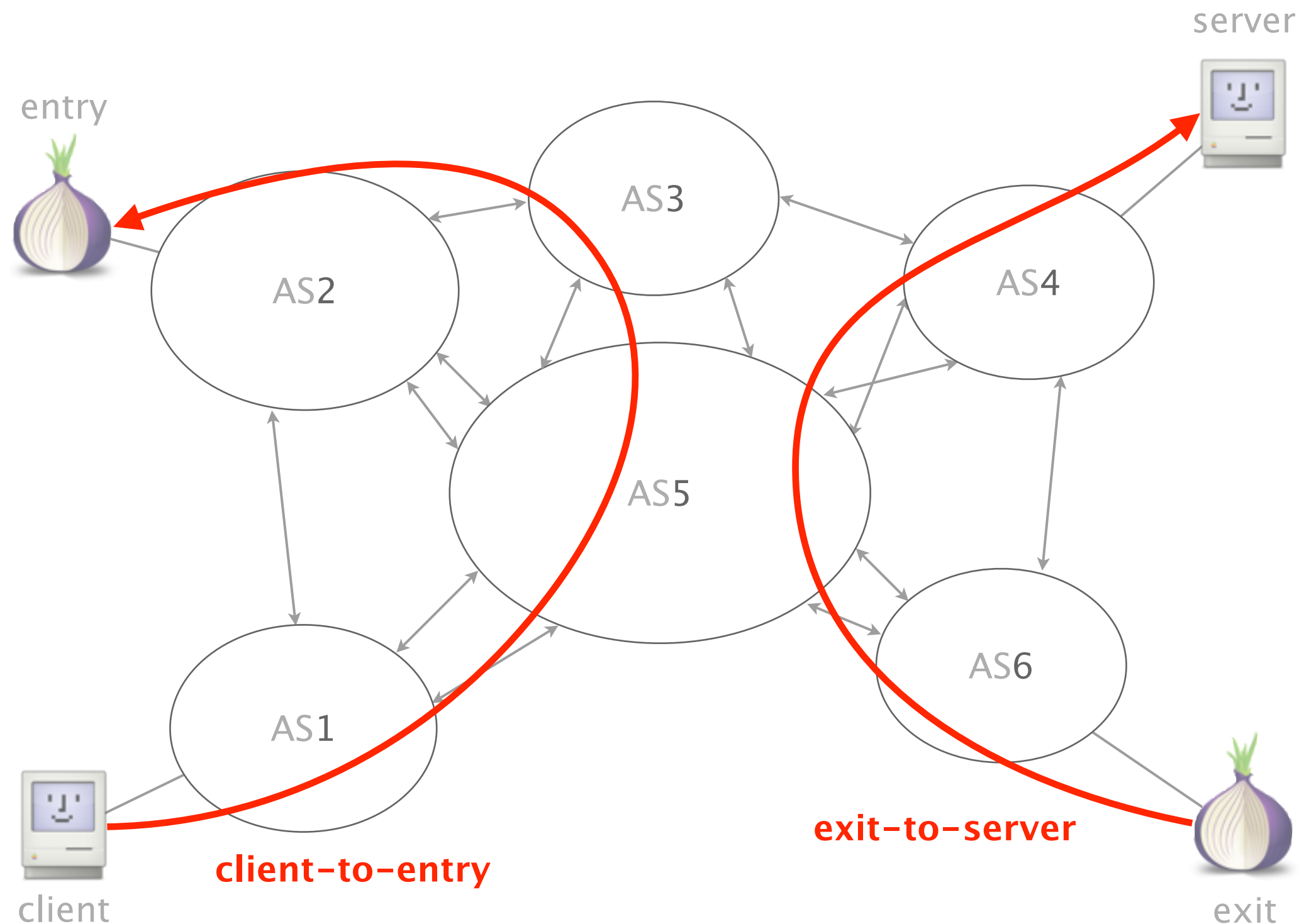
always sees traffic between
its customer and entry relays

During the attack, it also
saw traffic to/from exit relays
for a non-trivial fraction of traffic

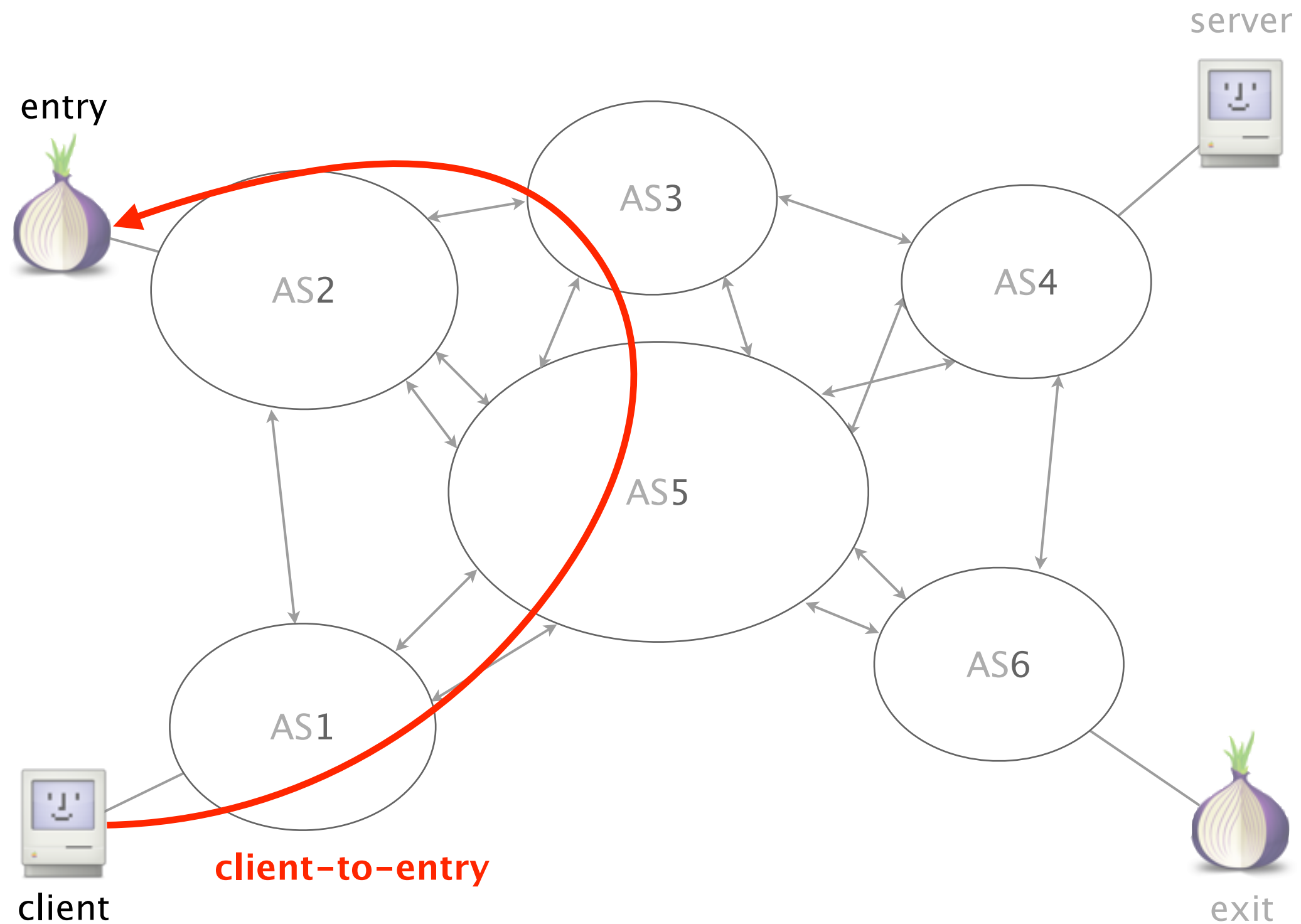
Intentional? No one knows.

Attack#3: Asymmetric routing, too,
increases the # of AS-level adversaries

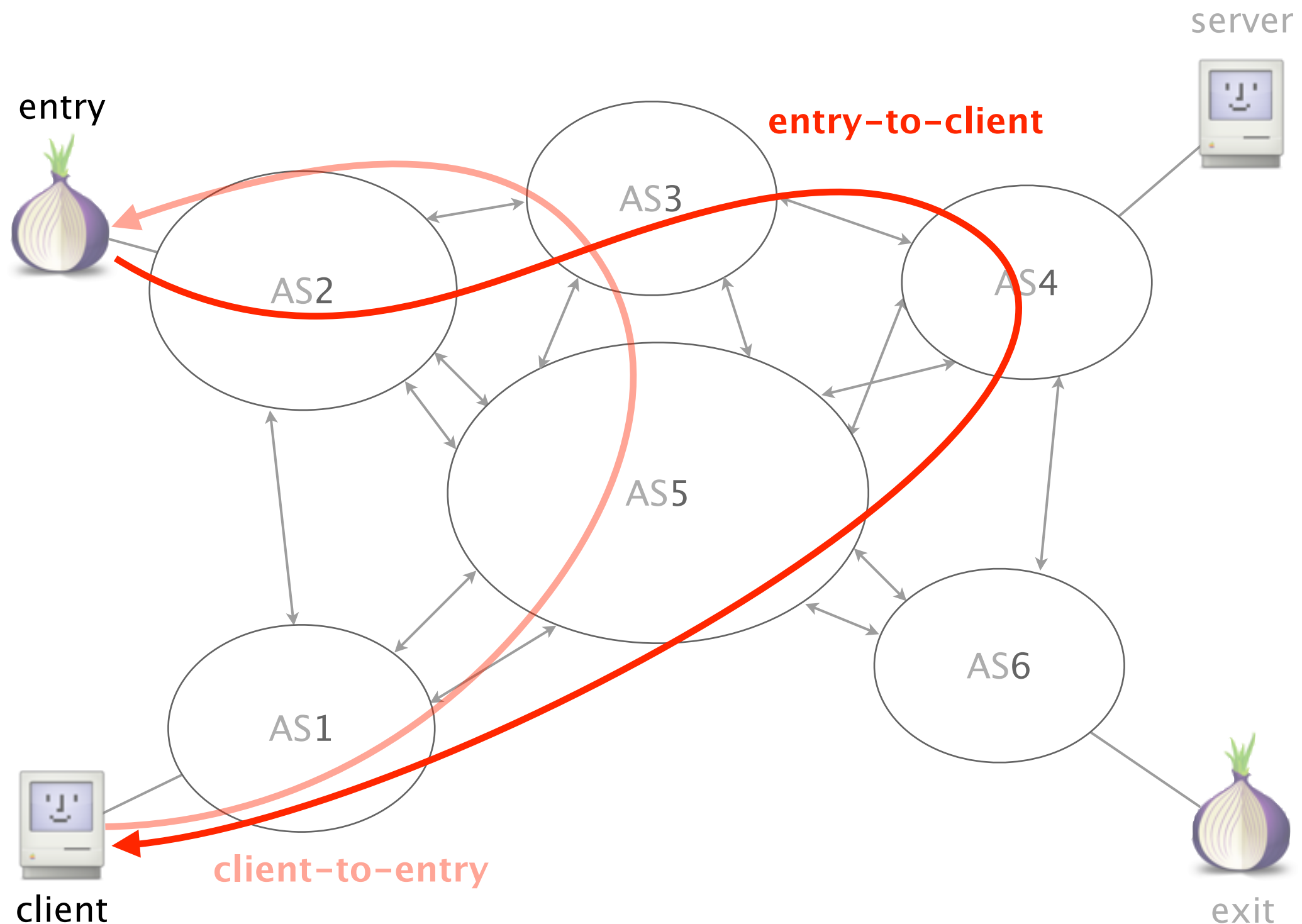
So far, we have considered one side of the Tor traffic:
client-to-entry and exit-to-server



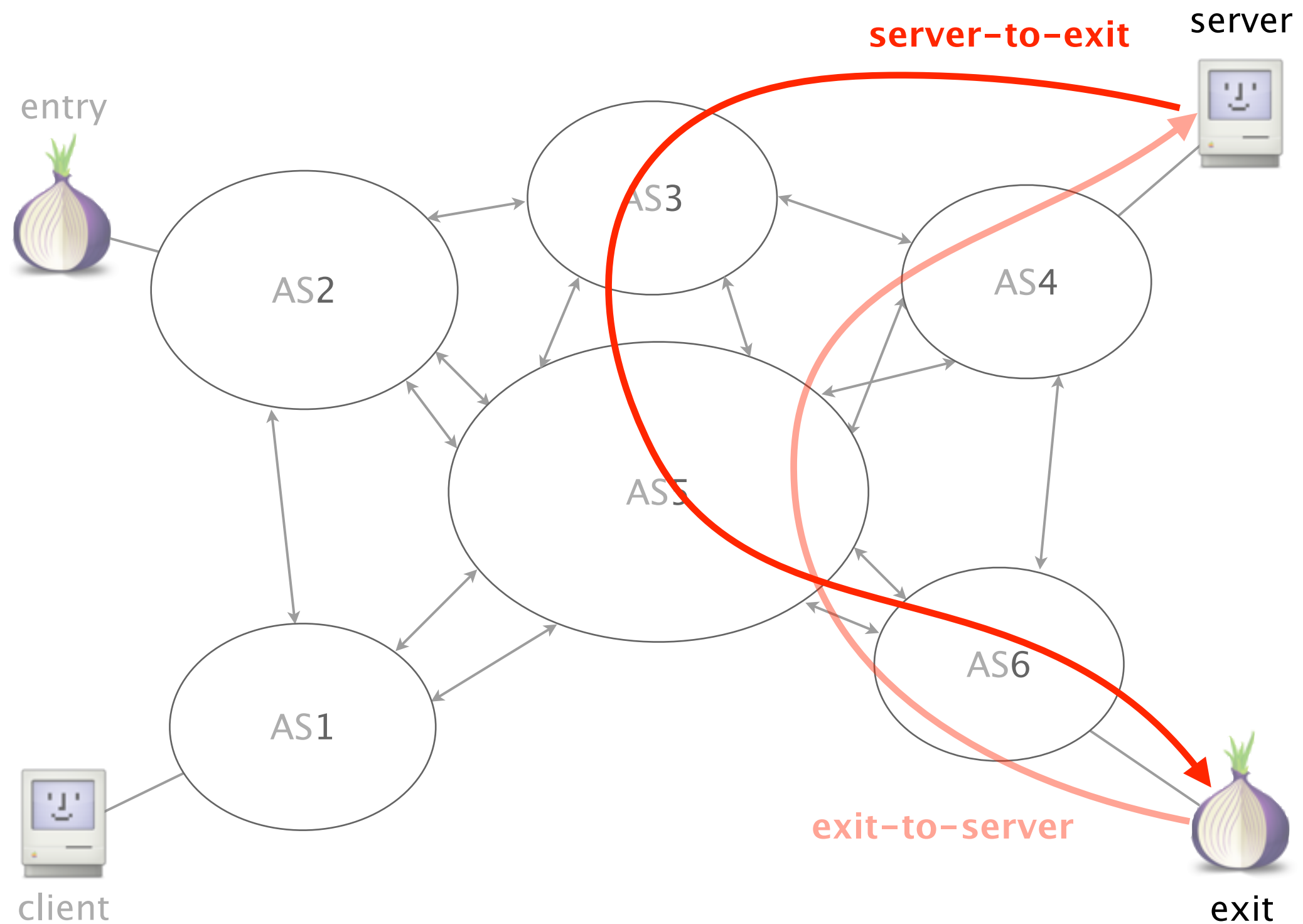
However, because of policies,
routing is often *asymmetric*



While AS4 does not see client-to-entry traffic, it sees entry-to-client traffic



Same applies for server-to-exit traffic



In terms of timing information,
both side of a TCP connection are highly correlated

In terms of timing properties,
both side of a TCP connection are highly correlated

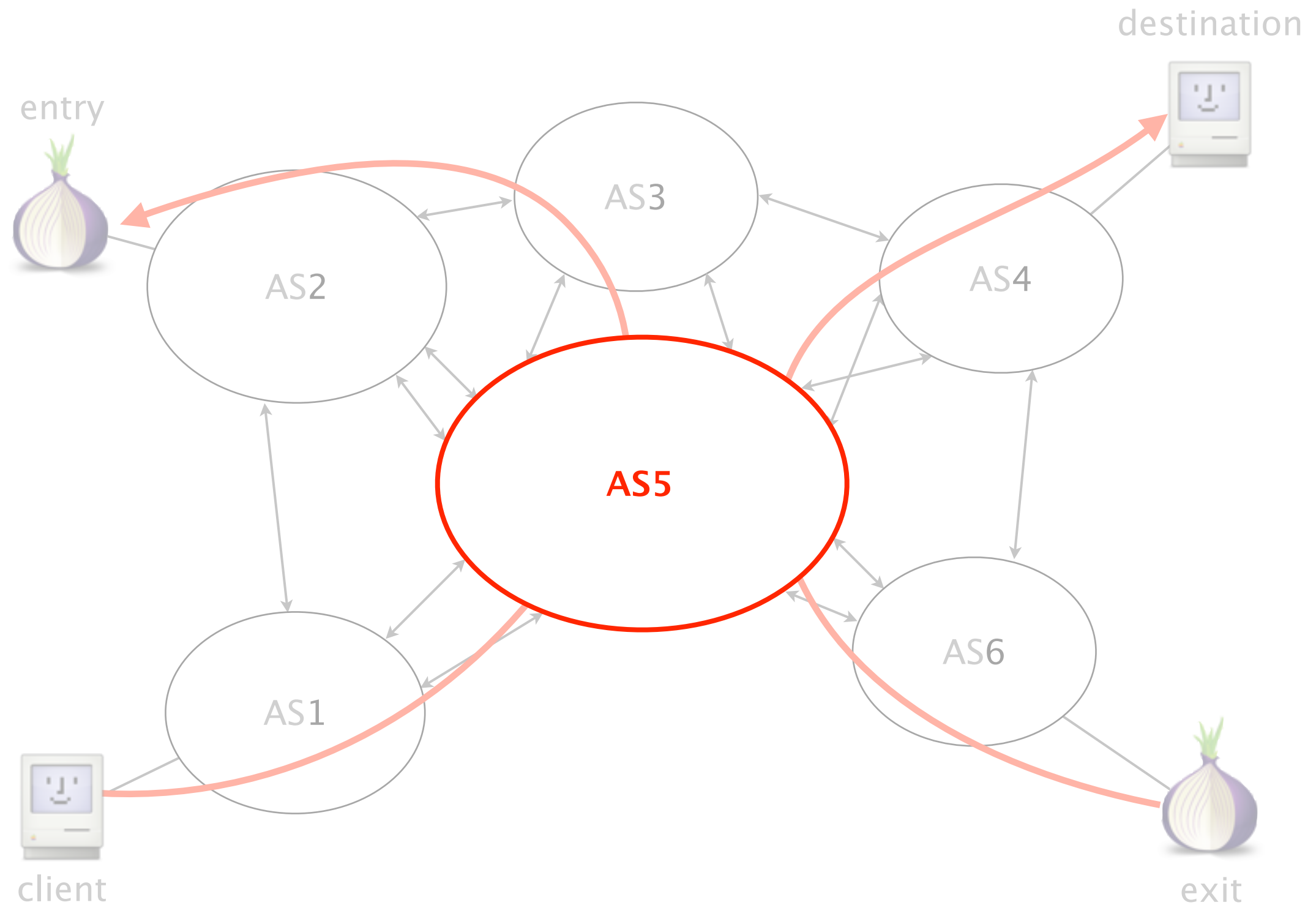
When collecting TCP
timing information,

seeing one direction
is almost equivalent to
seeing two directions

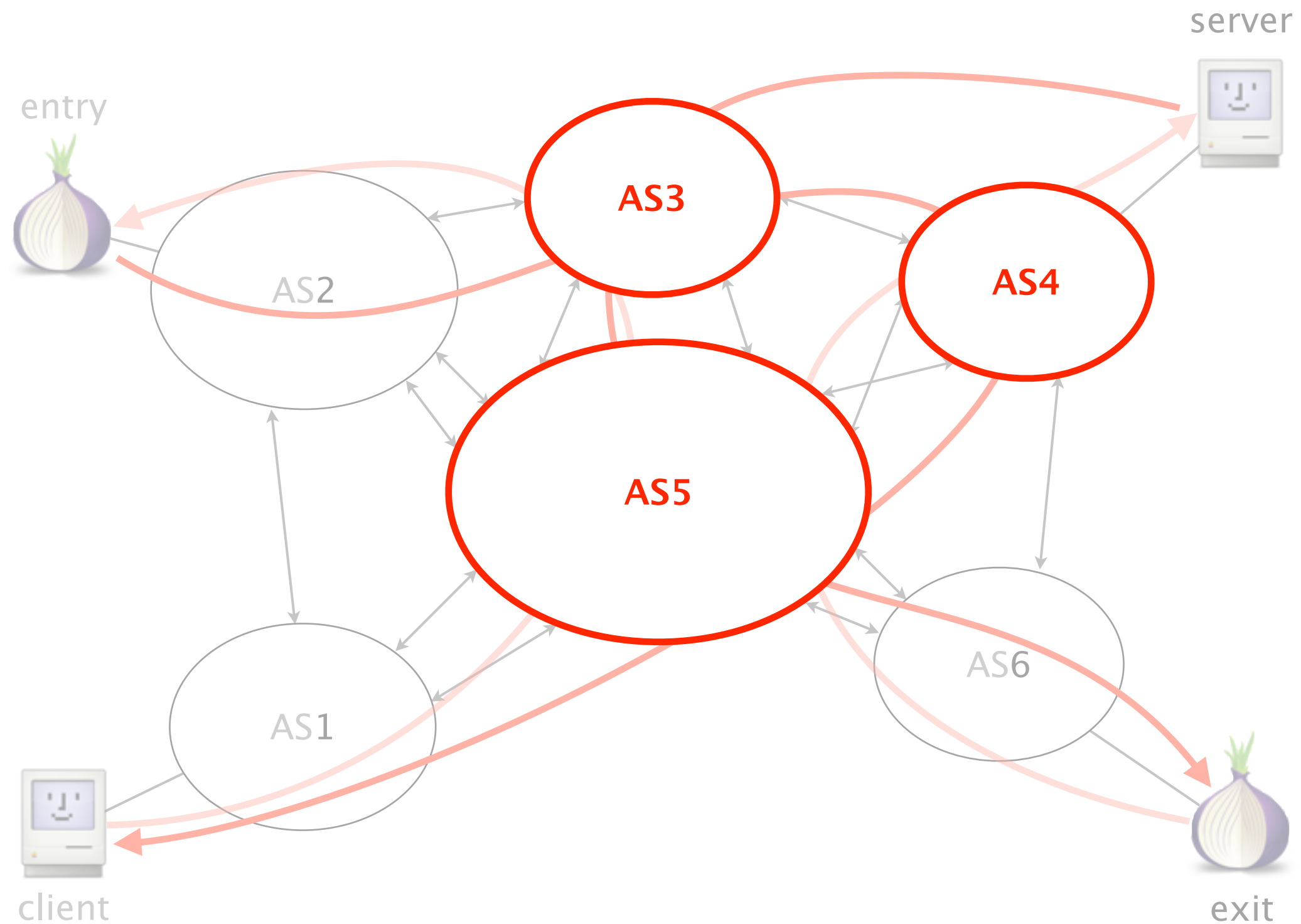
(*e.g.*, data packets)

(ACKs & data packets)

Considering only one direction,
1 AS is potentially compromising



Considering both directions,
3 ASes are potentially compromising



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Attacks

All your traffic belongs to me

2

Preliminary results

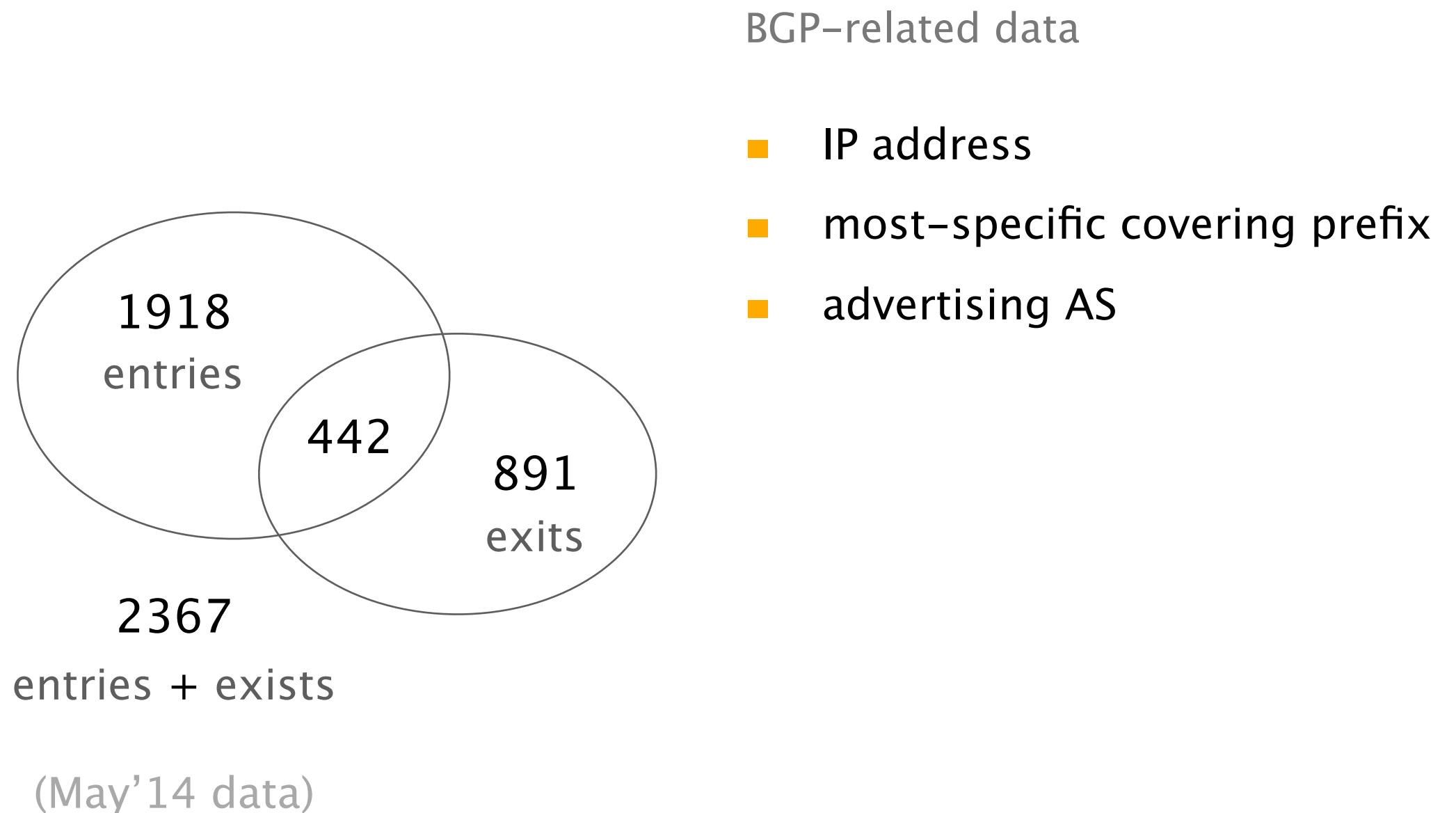
Eyes wide open

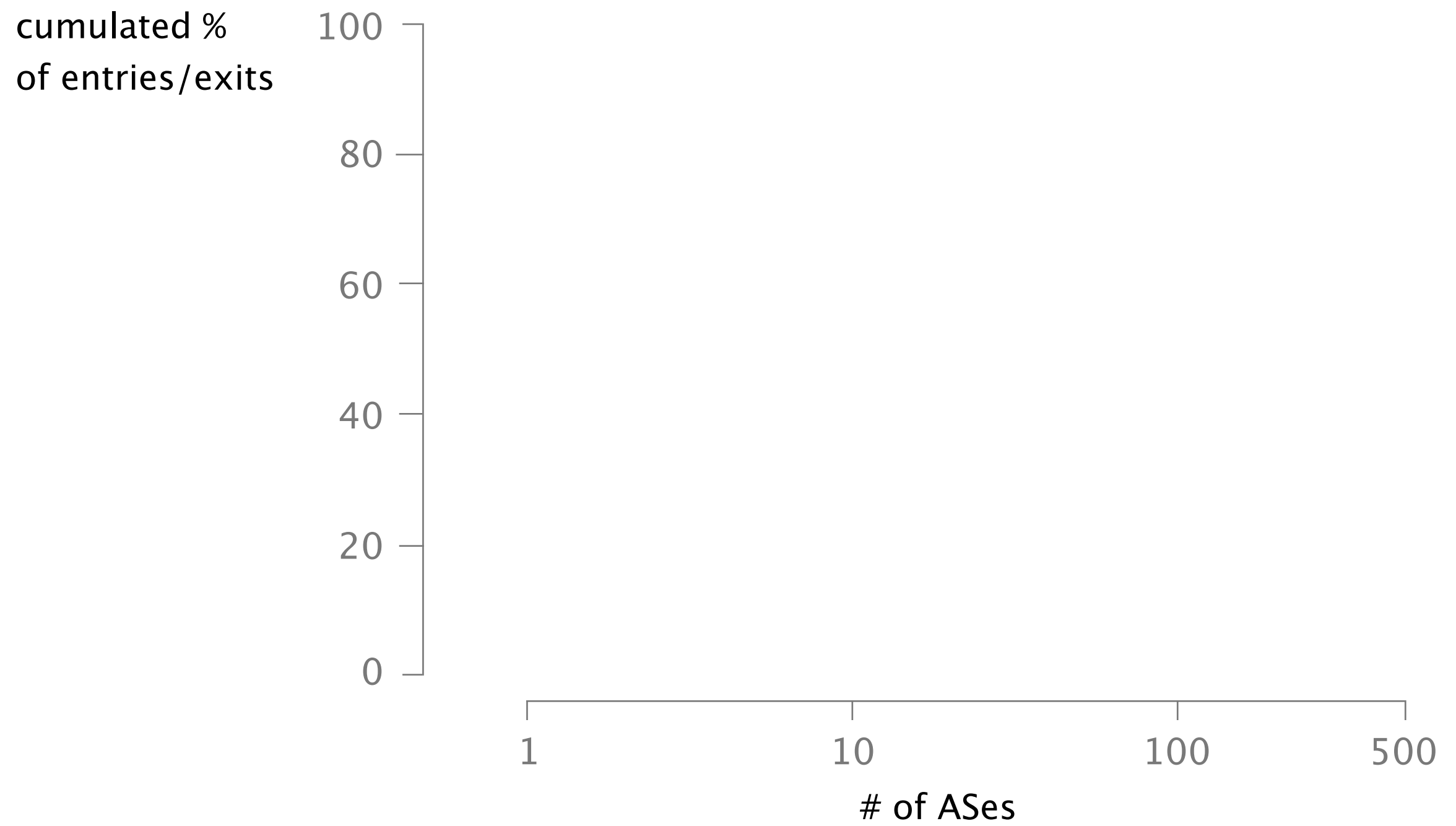
Countermeasures

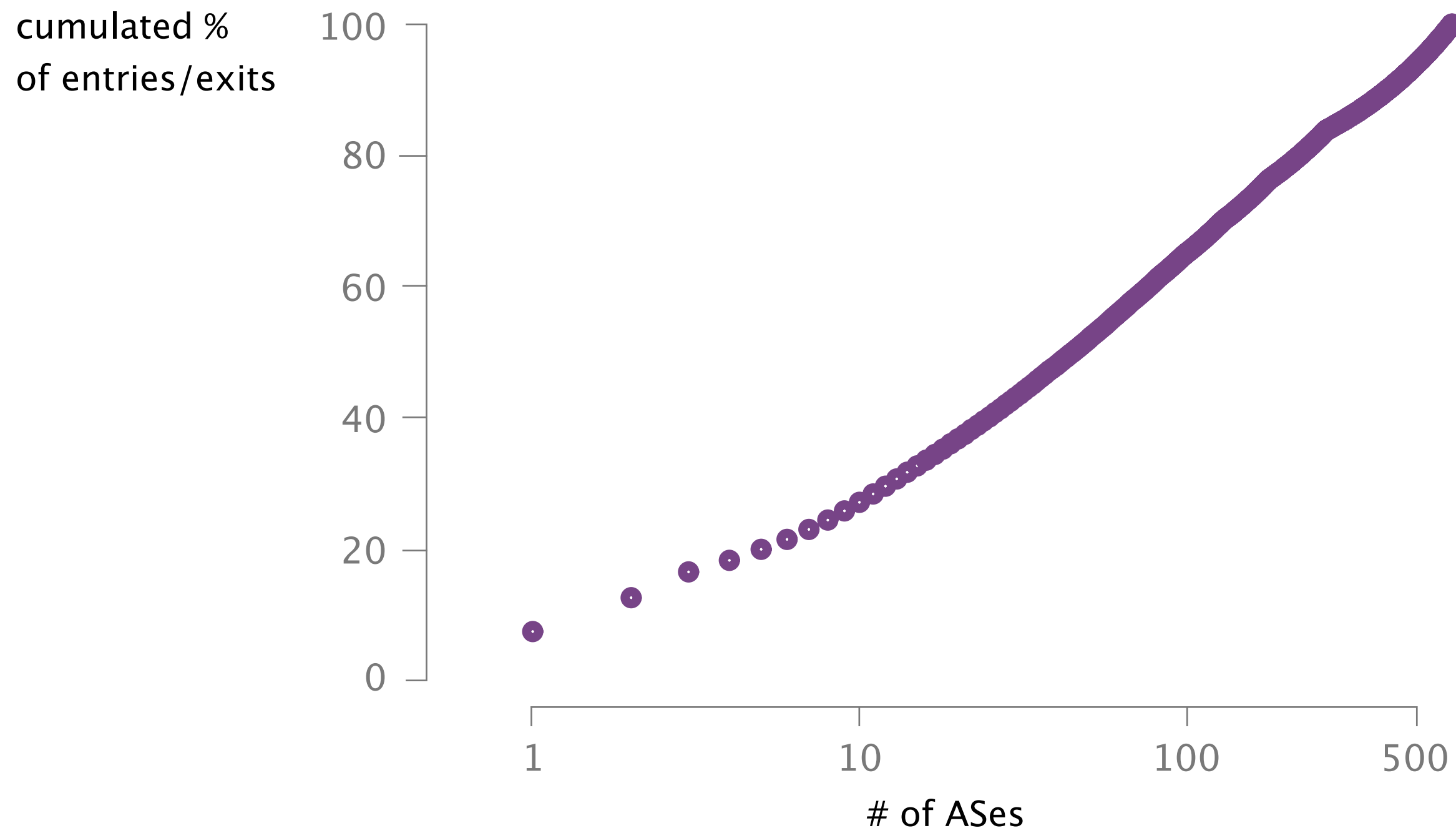
Close the curtains

Question#1: How many networks host entry and exit relays?

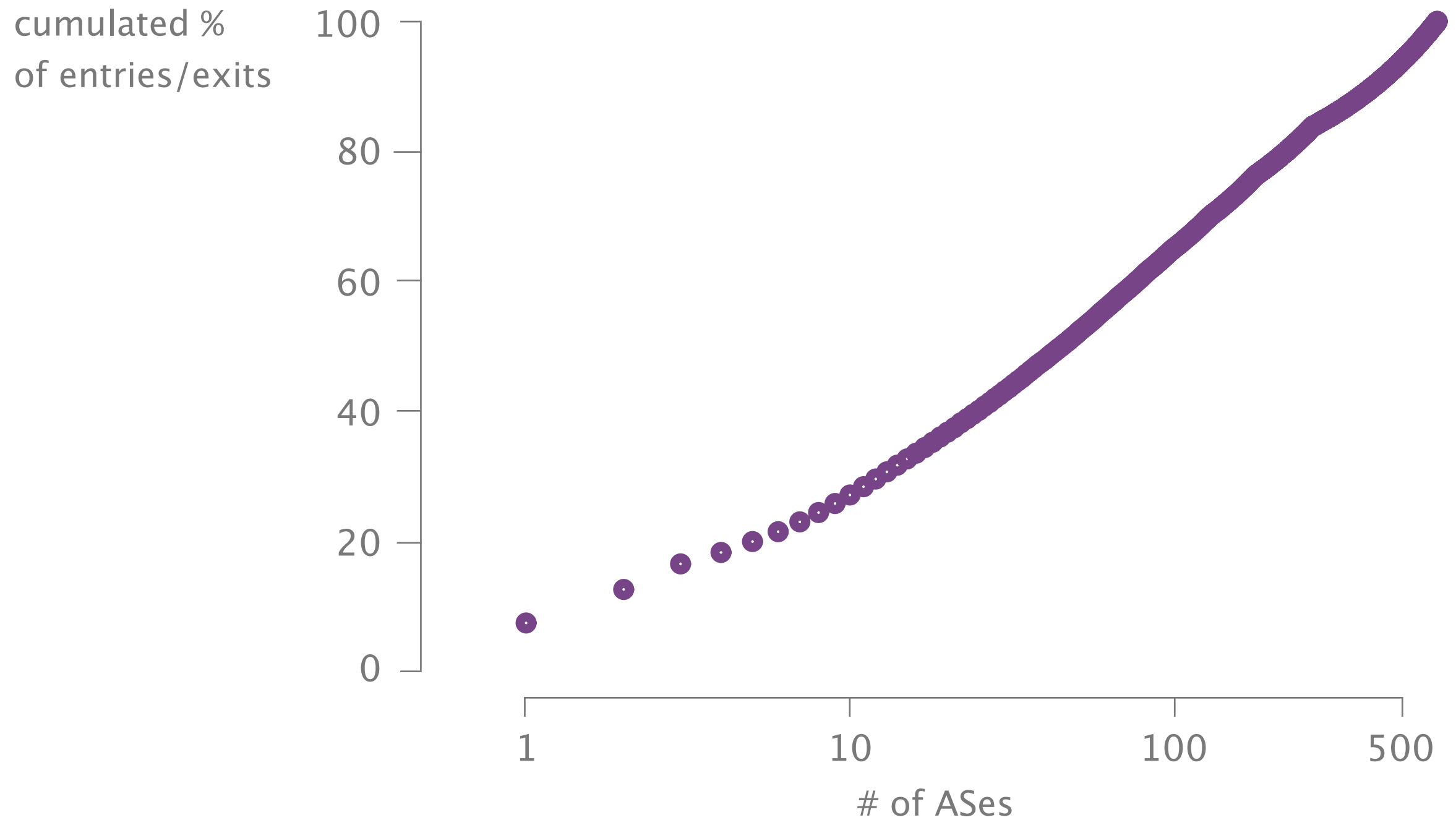
We collected BGP-related information
for each Tor entry and exit relay



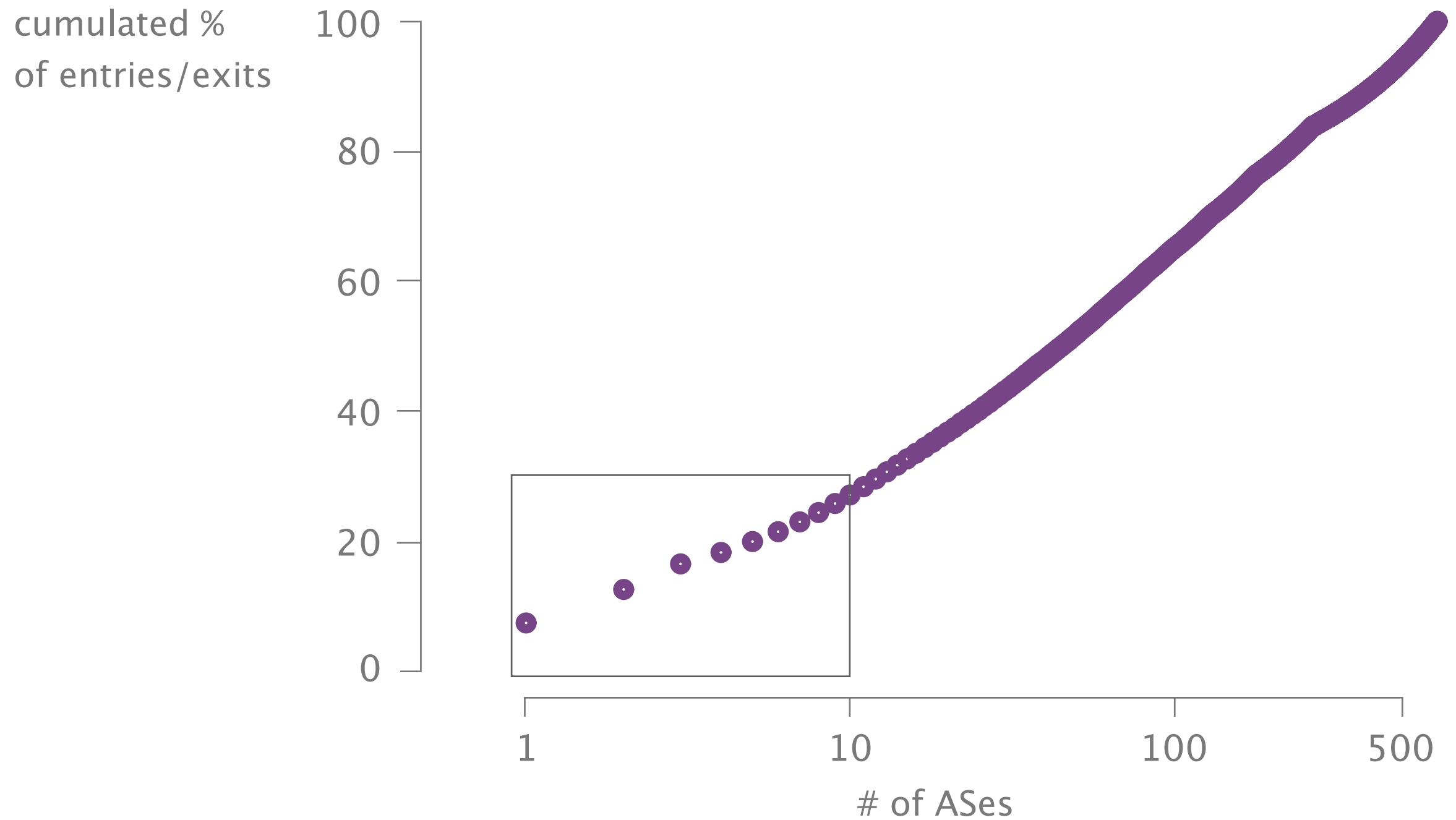




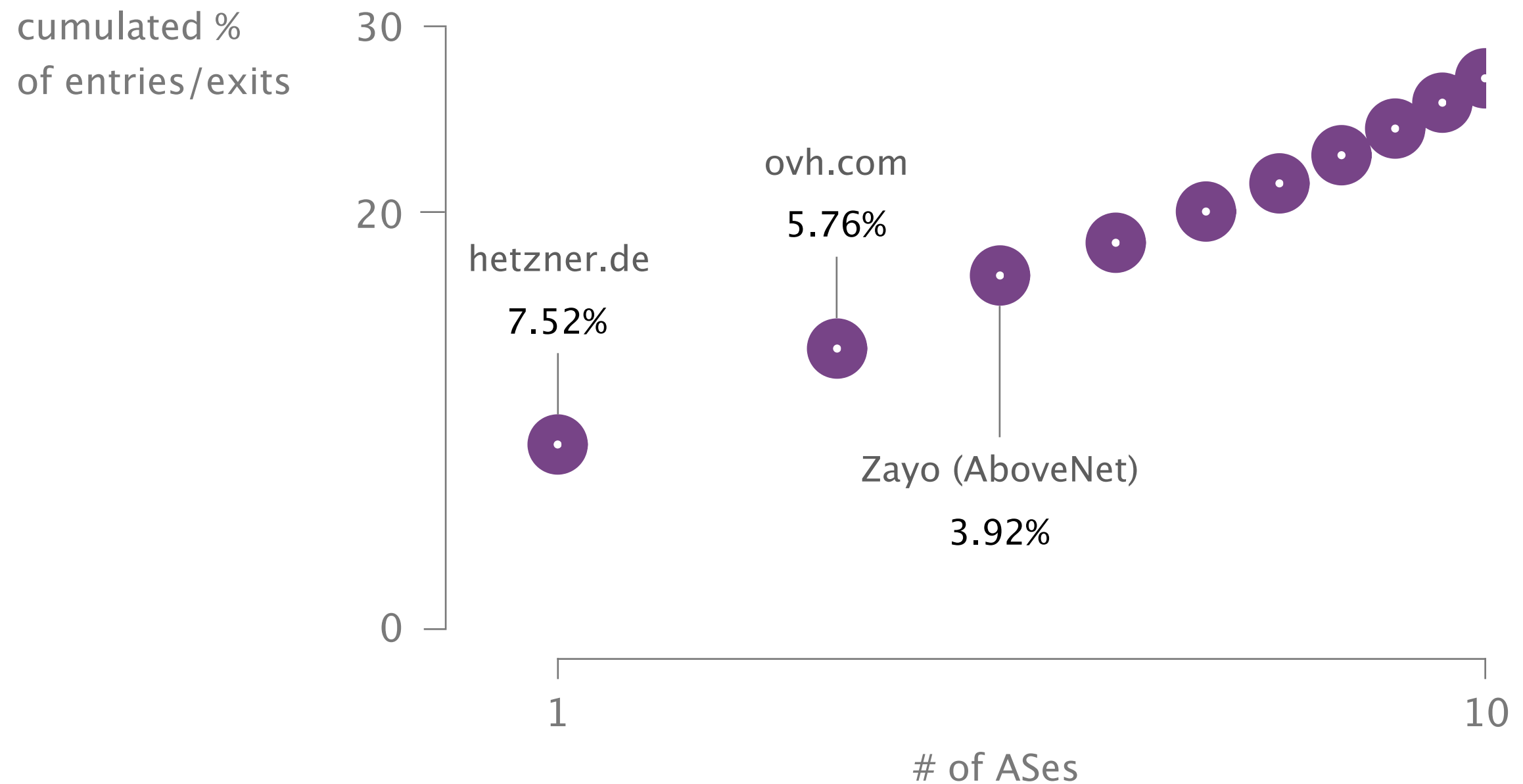
Entry & exit relays are concentrated in few ASes



Entry & exit relays are concentrated in few ASes



3 ASes host close to 20% of the entry & exit relays



Question#2: How much path changes were Tor prefixes seeing with respect to BGP prefixes?

To measure the effect of BGP dynamics
we collected BGP updates over 1 month

# BGP sessions (RIPE RIS collectors)	71
# BGP prefixes advertised by	1.2k 650 ASes
# BGP updates announcements/withdraws	1.4M

CCDF

100

80

60

40

20

0

.2

1

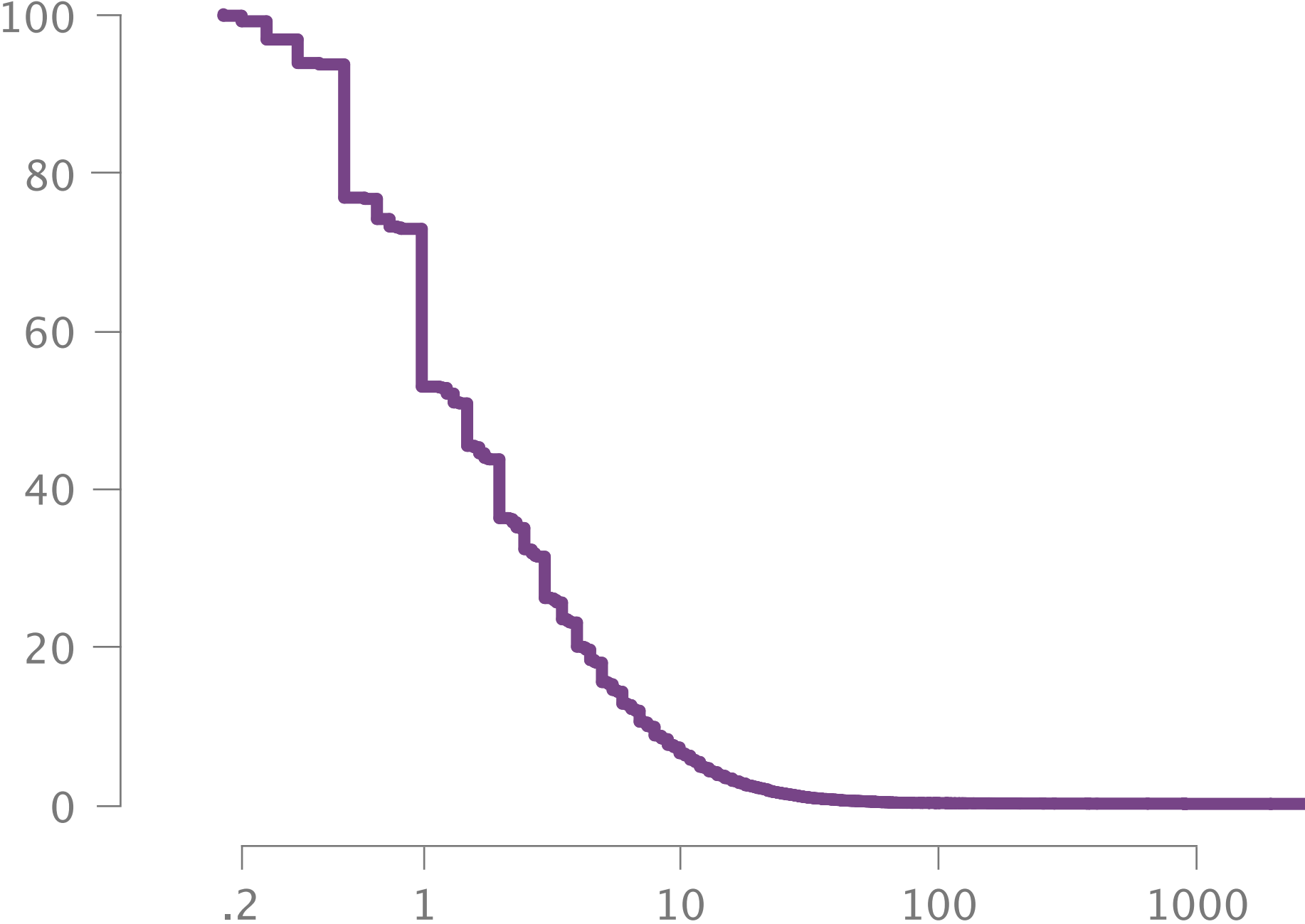
10

100

1000

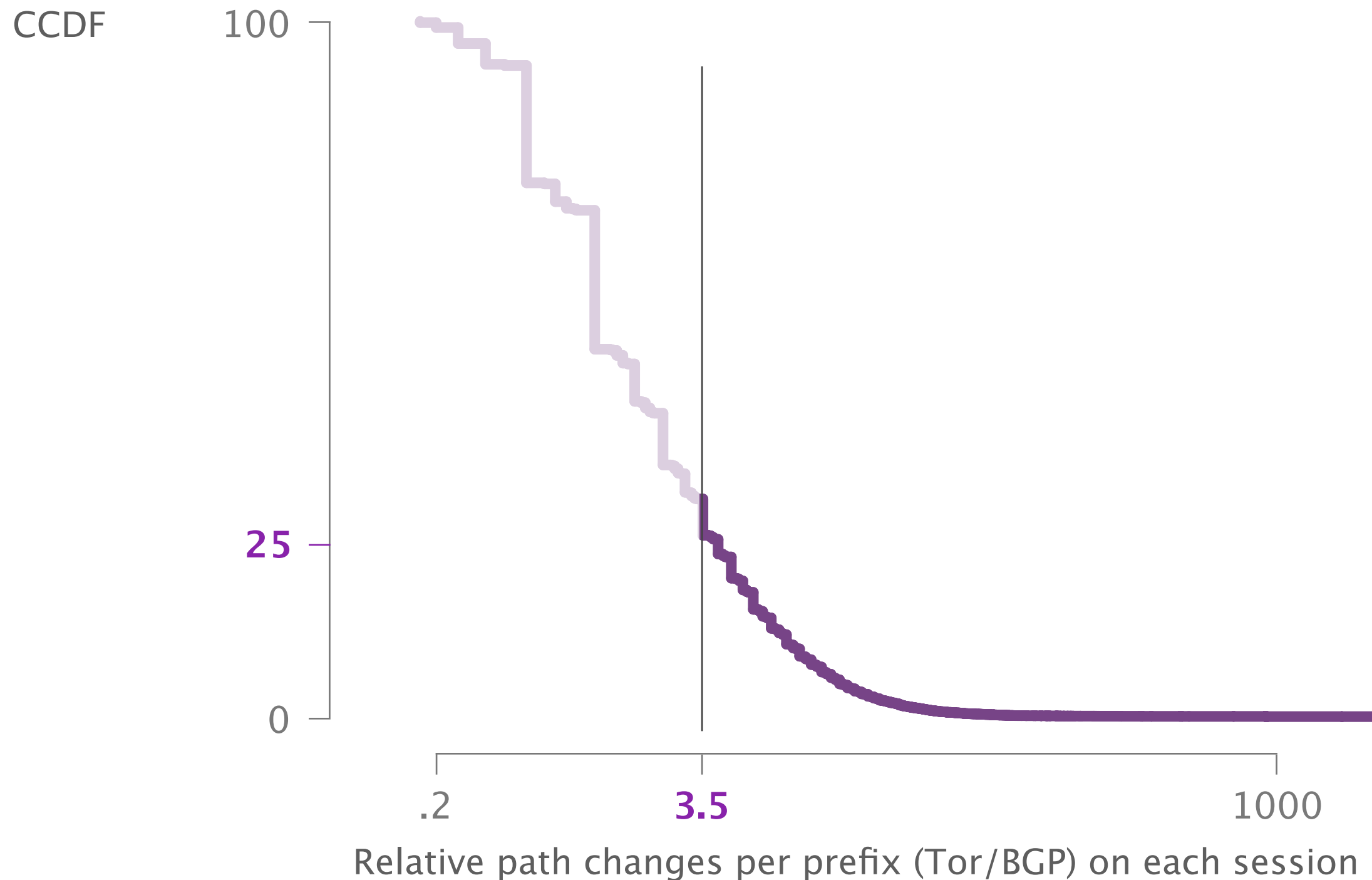
Relative path changes per prefix (Tor/BGP) on each session

CCDF



Relative path changes per prefix (Tor/BGP) on each session

In 25% of the cases, Tor prefixes saw >3.5 more changes than BGP prefixes on a session



These changes caused a bunch of extra ASes to see Tor traffic

In **60%** of the cases, **>2 extra ASes** receive traffic over the month because of BGP dynamics

significant as the average # of ASes per path is ~4

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Attacks

All your traffic belongs to me

Preliminary results

Eyes wide open

3

Countermeasures

Close the curtains

To protect itself, Tor should become
more aware of the network underlying it

Problems

Natural dynamism

Route manipulation

Asymmetric analysis

Countermeasures

prefer stable relays

discard “suspicious” relays
prefer close relays

encrypt transport header

Tools

BGP monitoring

BGP monitoring +
BGPsec

IPsec

These countermeasures help, but come with tradeoffs

Problems

Natural dynamism

Route manipulation

Asymmetric analysis

Countermeasures

prefer stable relays

discard “suspicious” relays
prefer close relays

encrypt transport header

Tradeoffs

more power to
fewer relays

not widely used
(easier to detect)

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Eyes wide open

Countermeasures

Close the curtains

BGP is not only a problem for Tor...

THREAT LEVEL

Hacker Redirects Traffic From 19 Internet Providers to Steal Bitcoins

BY ANDY GREENBERG 08.07.14 | 1:00 PM | [PERMALINK](#)

[f Share](#) 1.0k [T](#) Tweet 1,464 [g+1](#) 213 [in Share](#) 512 [Pin it](#)



... A bitcoin thief redirected a portion of online traffic from no less than 19 Internet service providers, including data from the networks of Amazon and other hosting services like DigitalOcean and OVH, with the goal of stealing cryptocurrency from a group of bitcoin users...

... A bitcoin thief redirected a portion of online traffic from no less than 19 Internet service providers, including data from the networks of Amazon and other hosting services like DigitalOcean and OVH, with the goal of stealing cryptocurrency from a group of bitcoin users...

OVH is the second AS in terms of # Tor relays hosted

... A bitcoin thief redirected a portion of online traffic from no less than 19 Internet service providers, including data from the networks of Amazon and other hosting services like DigitalOcean and **OVH**, with the goal of stealing cryptocurrency from a group of bitcoin users...

Internet routing matters when it comes to user anonymity

BGP dynamics decreases user anonymity over time
natural & induced, exacerbated by asymmetric routing

Initial results illustrate the vulnerabilities
full evaluation is required—and underway

Short-term countermeasures helps, to an extent
need a better understanding on their impacts

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