

# Seamless Network-Wide Migrations



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Network Automation 2011

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Human factors are responsible for  
50 to 80% of network device outages.

— Juniper Networks, 2008

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50 to 80% of network device outages.

— Juniper Networks, 2008

Large (resp. medium) businesses lose  
an average of 3.6% (resp. 1%) in annual  
revenue due to network downtime.

— Infonetics Research, 2004 – 2006

seamless  
migrations

The progressive modification of the  
configuration of a running network  
without loosing any IP packets

# Migrating the network can provide immediate benefits

Network migration can improve the

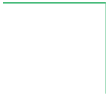
- manageability
- performance
- stability
- security

of the entire network

# Migrating the network can provide immediate benefits

Network migrations can improve the

- manageability
- performance
- stability
- security



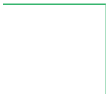
- introduce (remove) a hierarchy
- support different features
- improve route manipulation

of the entire network

# Migrating the network can provide immediate benefits

Network migrations can improve the

- manageability
- **performance**
- stability
- security




- reduce control-plane load
- reduce routing-table size
- reduce convergence time

of the entire network

# Migrating the network can provide immediate benefits

Network migrations can improve the

- manageability
- performance
- **stability**
- security



isolate parts of the network  
reduce the churn

of the entire network




# Migrating the network can provide immediate benefits

Network migrations can improve the

- manageability
- performance
- stability
- security

of the entire network



use stable code base  
prevent potential attacks

# Migrating the network is operationally complex

Reconfigure a running network  
while respecting Service Level Agreement

Make highly distributed changes  
on all the routers, in a coordinated manner

Face potential routing anomalies  
as non-migrated routers interact with migrated ones

# Seamless Network-Wide Migrations



- 1 IGP migrations  
Ordering matters
- 2 BGP migrations  
Ongoing work
- 3 How can we help ?  
Manage complexity

# Seamless Network-Wide Migrations



1

IGP migrations

Ordering matters

BGP migrations

Ongoing work

How can we help ?

Manage complexity

Problem                Replace the IGP configuration of a  
running network, router-by-router,  
without causing any anomalies

Details in:        Laurent Vanbever, Stefano Vissicchio, Cristel Pelsser, Pierre Francois and Olivier Bonaventure.  
*Seamless Network-Wide IGP Migrations*. ACM SIGCOMM 2011, Aug. 2011.

Sub-problem 1

Replace the IGP configuration of a running network, router-by-router, without causing any anomalies

Current Practice

Run the two IGP configurations in parallel

# Migrating the IGP usually requires running two routing planes

Abstract model of a router

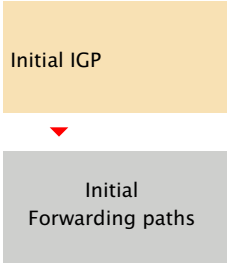
Control-plane

Initial IGP

At first, the initial IGP dictates the forwarding paths being used

Data-plane

Initial  
Forwarding paths



# Migrating the IGP usually requires running two routing planes

Abstract model of a router

Control-plane

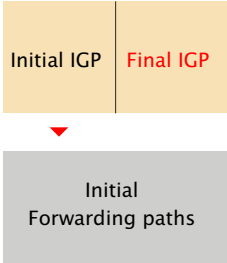
Initial IGP

Final IGP

Then, the final IGP is introduced without changing the forwarding

Data-plane

Initial  
Forwarding paths



The diagram illustrates the migration of the IGP in a router. It is divided into two main sections: Control-plane and Data-plane. The Control-plane is further divided into Initial IGP and Final IGP. The Data-plane contains Initial Forwarding paths. A red arrow points from the Control-plane to the Data-plane, indicating the flow of information or the migration process.



# Migrating the IGP usually requires running two routing planes

Abstract model of a router

Control-plane

Initial IGP

Final IGP

After having converged, the final IGP is used by flipping the preference

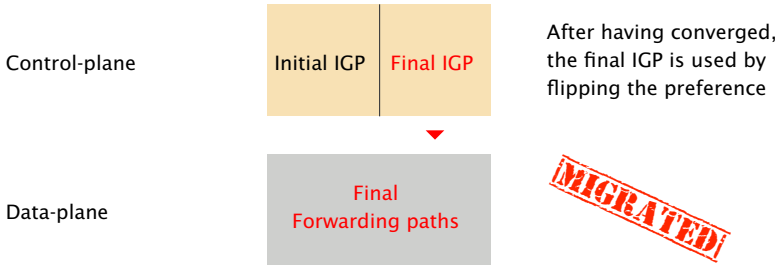
Data-plane

Final  
Forwarding paths



# Migrating the IGP usually requires running two routing planes

Abstract model of a router



# Migrating the IGP usually requires running two routing planes

Abstract model of a router

Control-plane

Final IGP

The initial IGP is removed as it is not used anymore

Data-plane

Final  
Forwarding paths

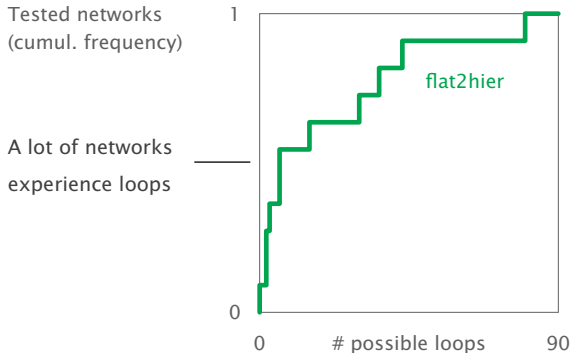
**MIGRATED**

Sub-problem 1

Replace the IGP configuration of a running network, router-by-router, without causing any anomalies

Sub-problem 2      Replace the IGP configuration of a running network, router-by-router, without causing any anomalies

# Migrating the IGP can create *migration loops*



Up to 90 *migration loops* can arise during an IGP migration

Sub-problem 2      Replace the IGP configuration of a running network, router-by-router, without causing any anomalies

Sub-problem 2      Replace the IGP configuration of a running network, router-by-router, without causing any anomalies

Contributions      Seamless IGP migration is possible as long as the reconfiguration process follows a strict ordering



## Contributions

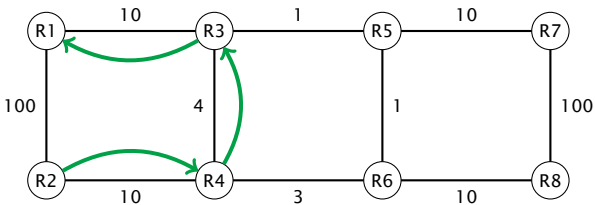
Seamless IGP migration is possible as long as the reconfiguration process follows a **strict ordering**

|  
which one ?

## Reconfiguring the IGP might change the forwarding paths being used

In a flat IGP, routers forward traffic according to the shortest-path towards the destination.

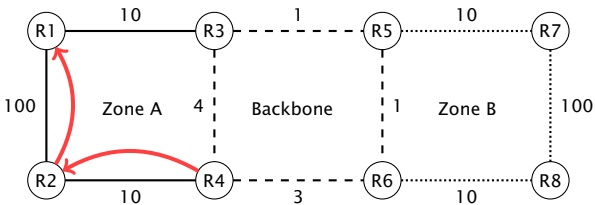
In a flat IGP, R2 reaches R1 via R4



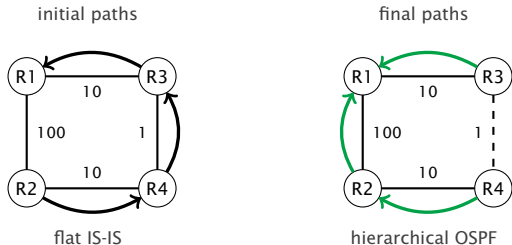
## Reconfiguring the IGP might change the forwarding paths being used

In a hierarchical IGP, routers prefer paths contained within a single zone over the ones crossing several zones

In a hierarchical IGP, R2 reaches R1 directly

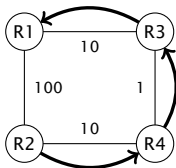


Whenever the forwarding paths change,  
forwarding loops can be created



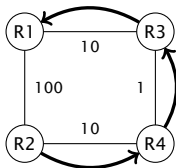
Forwarding paths towards R1

initial paths

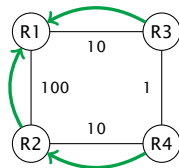


flat IS-IS

intermediate paths



final paths

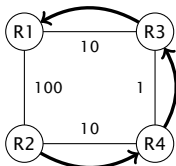


hierarchical OSPF

Forwarding paths towards R1

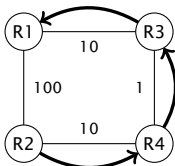
## First, we migrate R3

initial paths

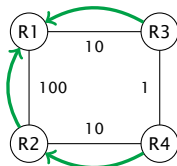


flat IS-IS

intermediate paths



final paths

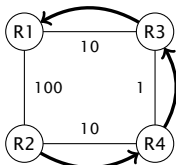


hierarchical OSPF

Forwarding paths towards R1

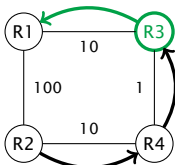
First, we migrate R3

initial paths

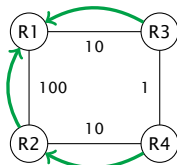


flat IS-IS

intermediate paths



final paths

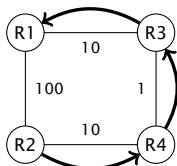


hierarchical OSPF

Forwarding paths towards R1

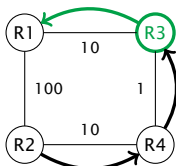
Then, we migrate R4

initial paths

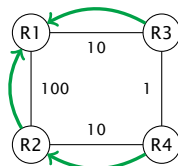


flat IS-IS

intermediate paths



final paths



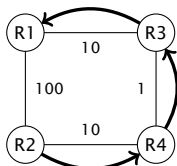
hierarchical OSPF

Forwarding paths towards R1



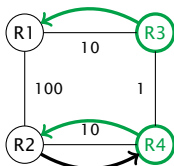
Then, we migrate R4

initial paths



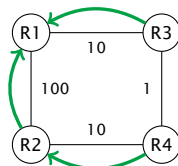
flat IS-IS

intermediate paths



Forwarding paths towards R1

final paths

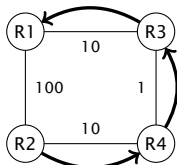


hierarchical OSPF

Whenever the forwarding paths change,  
forwarding loops can be created

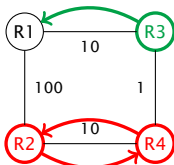
A loop is created if R4 is migrated before R2

initial paths

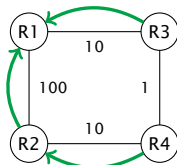


flat IS-IS

intermediate paths



final paths



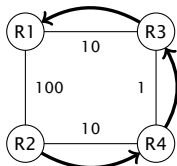
hierarchical OSPF

Forwarding paths towards R1

# Migrations have to be performed following a precise ordering

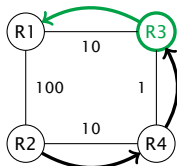
No loop arises if R2 is migrated before R4

initial paths

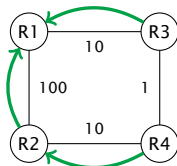


flat IS-IS

intermediate paths



final paths



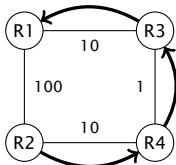
hierarchical OSPF

Forwarding paths towards R1

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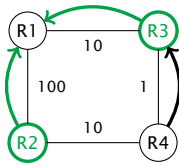
No loop arises if R2 is migrated before R4

initial paths

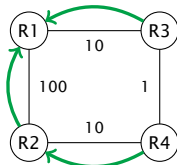


flat IS-IS

intermediate paths



final paths



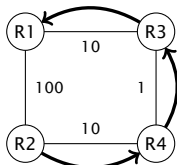
hierarchical OSPF

Forwarding paths towards R1

# Migrations have to be performed following a precise ordering

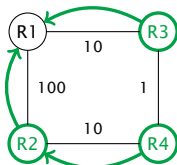
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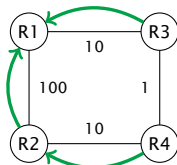


flat IS-IS

intermediate paths



final paths



hierarchical OSPF

Forwarding paths towards R1

Although hard in theory, finding an ordering is possible in practice

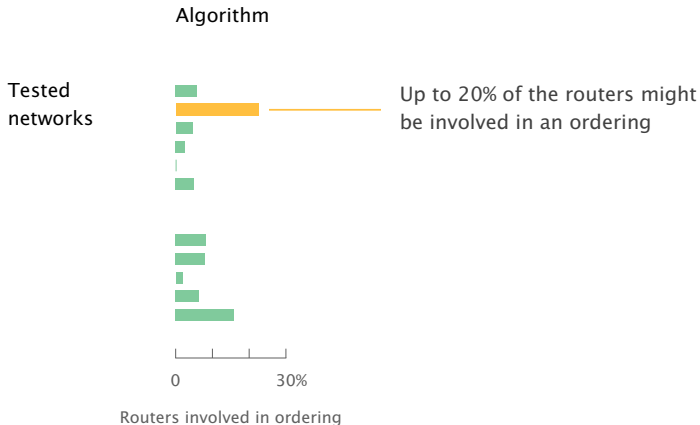
Implementation of two ordering algorithms

1. correct, complete, but slow
2. correct, not-complete, but fast

Complete support of the

- introduction (removal) of an IGP hierarchy
- introduction (removal) of route summarization
- modifications of link weights

For all the tested networks,  
we have been able to find an ordering



Our techniques can also deal with links and nodes failures during the migration

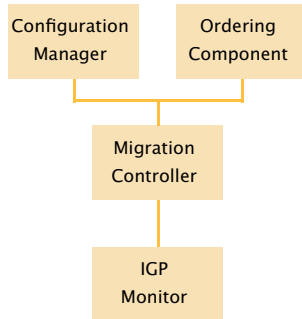
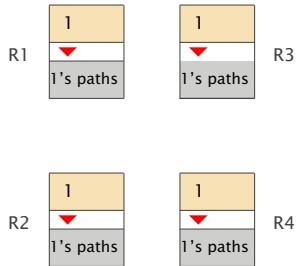
Failures can change the computed ordering as they modify the underlying IGP topology

Solutions

- Precompute failover orderings
- Compute a new ordering when a failure is detected

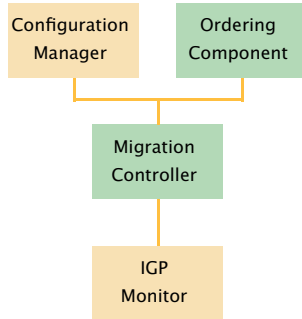
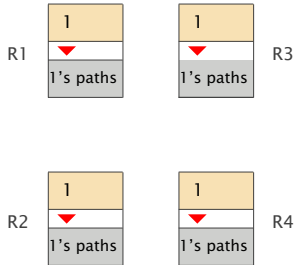


We implemented a provisioning system which automates the process



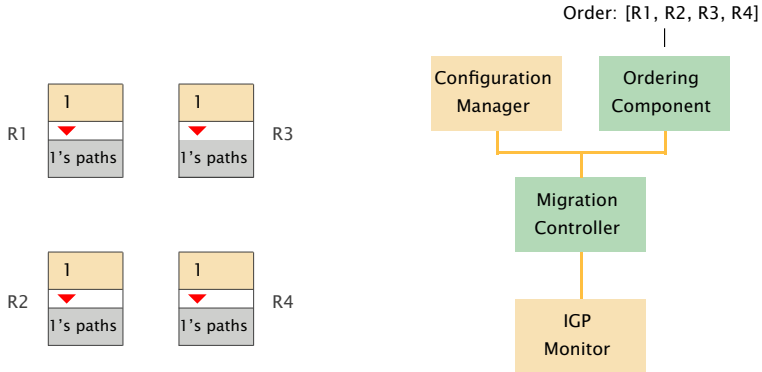
Network in which IGP 1 is replaced by IGP 2

First, the *Ordering Component* computes the ordering (if any)



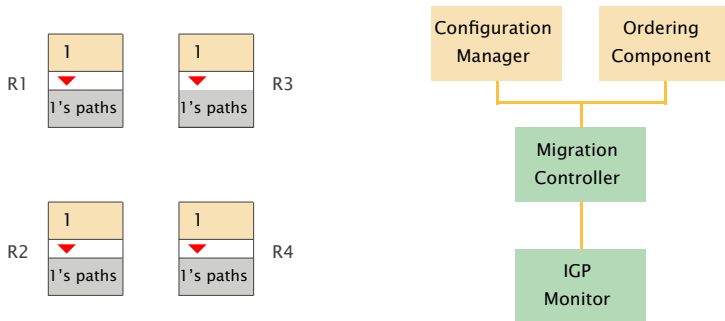
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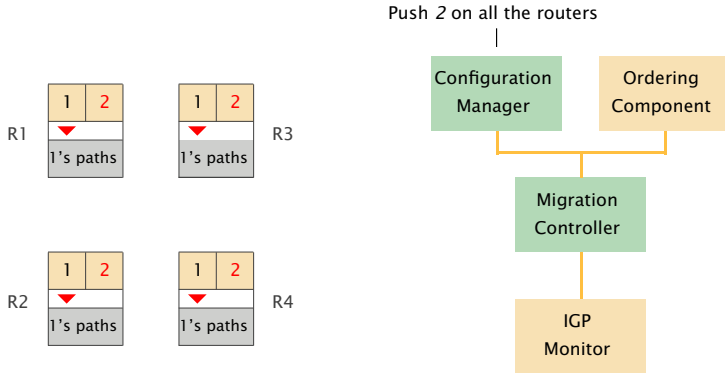
Network in which IGP 1 is replaced by IGP 2

Second, the *IGP Monitor* builds a dynamic view of the IGP and assesses its stability



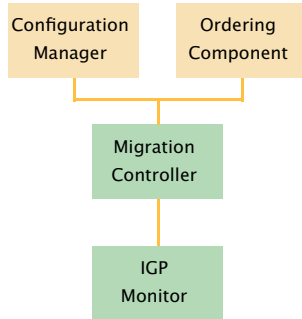
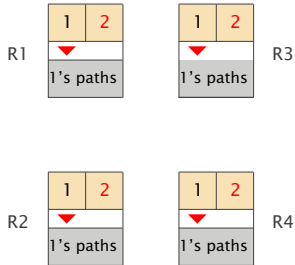
Network in which IGP 1 is replaced by IGP 2

Third, the *Configuration Manager* introduces the, final configuration (not yet used) on all the routers



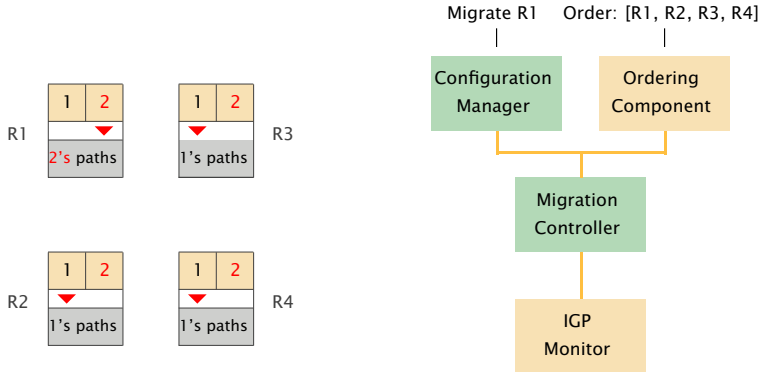
Network in which IGP 1 is replaced by IGP 2

Fourth, the final IGP's completeness and stability are verified by the *IGP Monitor*



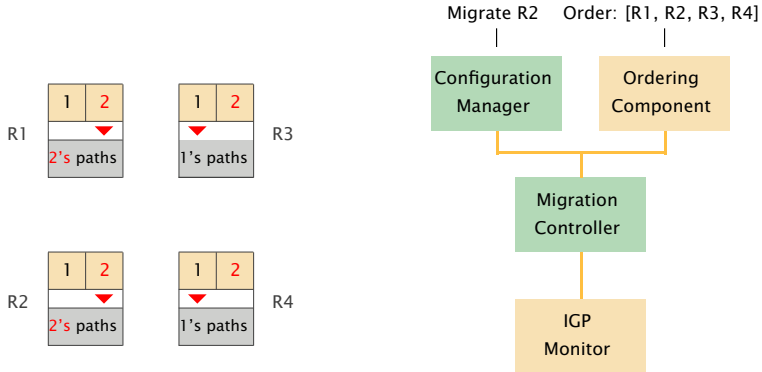
Network in which IGP 1 is replaced by IGP 2

Fifth, the *Configuration Manager* reconfigures each router – according to the ordering — so that it uses the final IGP



Network in which IGP 1 is replaced by IGP 2

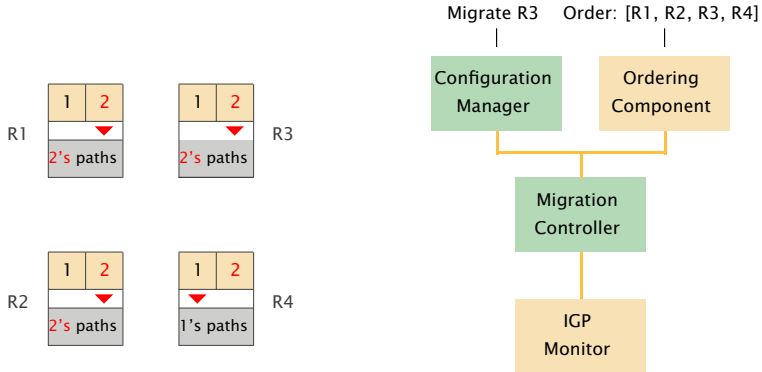
Fifth, the *Configuration Manager* reconfigures each router – according to the ordering — so that it uses the final IGP



Network in which IGP 1 is replaced by IGP 2

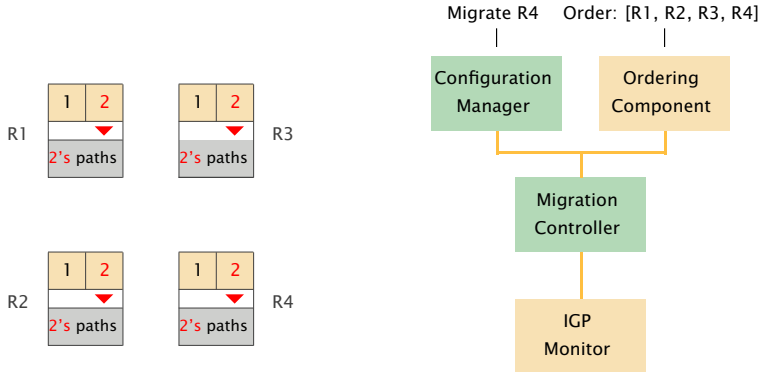


Fifth, the *Configuration Manager* reconfigures each router – according to the ordering — so that it uses the final IGP



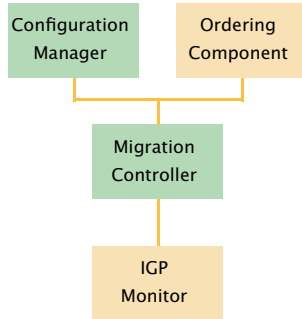
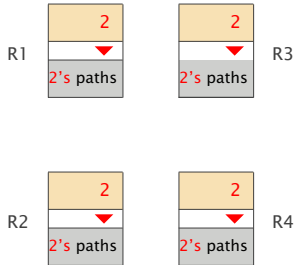
Network in which IGP 1 is replaced by IGP 2

Fifth, the *Configuration Manager* reconfigures each router – according to the ordering — so that it uses the final IGP



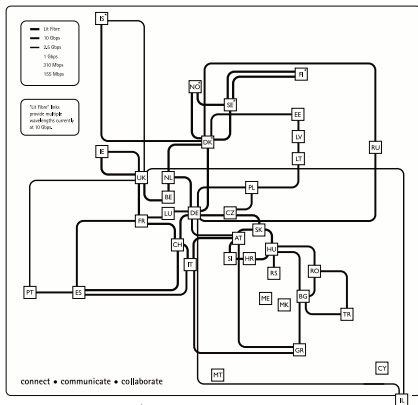
Network in which IGP 1 is replaced by IGP 2

Sixth, the IGP migration is over. The *Configuration Manager* removes the initial IGP configuration from each router



Network in which IGP 1 is replaced by IGP 2

# Let's reconfigure an existing network from a *flat* IGP ...



Planned Backbone Topology by the end of 2010. GEANT is operated by DANTE on behalf of Europe's NRENs.

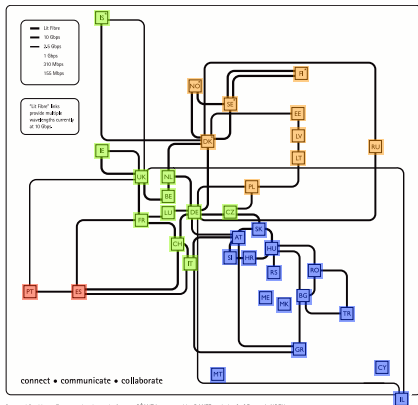
GEANT

European research network

36 routers

53 links

# Let's reconfigure an existing network from a *flat* IGP to a *hierarchical* IGP



Planned Backbone Topology by the end of 2010. GEANT is operated by DANTE on behalf of Europe's NRENs.

GEANT

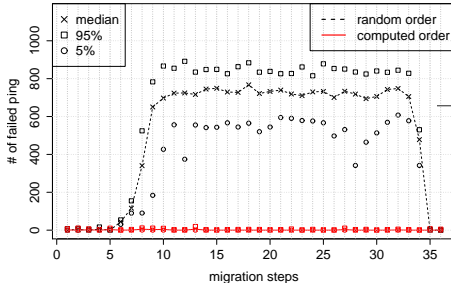
European research network

36 routers

53 links

- Backbone zone
- South-west zone
- South-east zone
- North-east zone

# Lossless reconfiguration is possible, by following the precomputed ordering



Traffic gets lost during  
more than 80% of the process

No loss occurs  
with proper ordering

Average results (50 repetitions) computed on 700+ pings  
per step from every router to 5 problematic destinations

# Seamless Network-Wide Migrations



IGP migrations

Ordering matters

2

**BGP migrations**

Ongoing work

How can we help ?

Manage complexity

# BGP reconfiguration scenarios are numerous

## Scenarios

iBGP

Full-mesh to route-reflection

Route-reflection to full-mesh

Add, (re-)move sessions

eBGP

Add, (re-)move sessions

Modify in(out)bound policies



# BGP reconfiguration scenarios are numerous ... and problematic

	Scenarios	Problems
iBGP	Full-mesh to route-reflection Route-reflection to full-mesh Add, (re-)move sessions	Forwarding loops Routing instabilities Visibility issues
eBGP	Add, (re-)move sessions Modify in(out)bound policies	Traffic shifts

# BGP reconfiguration scenarios are numerous ... and problematic

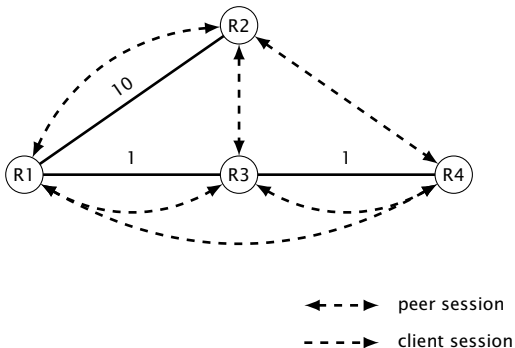
Preliminary evaluations confirm  
that problems may arise even when  
following the best current practices

## Problems

- Forwarding loops
- Routing instabilities
- Visibility issues
- Traffic shifts

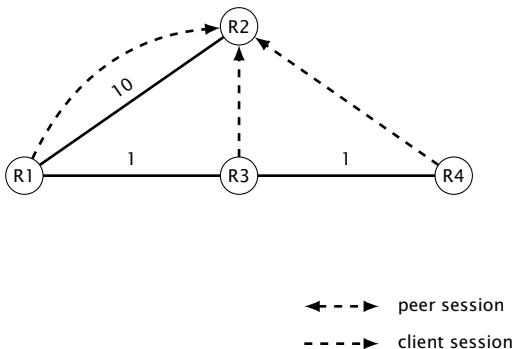
As for the IGP, reconfiguring BGP can create problems

Initial configuration



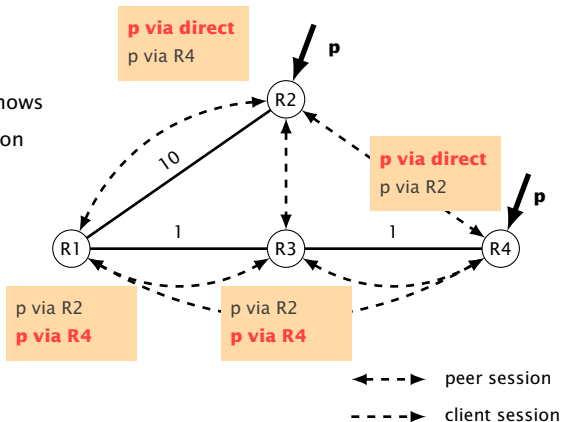
As for the IGP, reconfiguring BGP  
can create problems

Final configuration



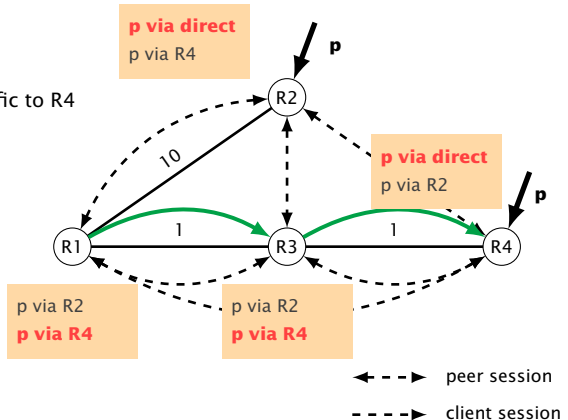
## As for the IGP, reconfiguring BGP can create problems

In a full-mesh, every router knows  
2 paths to reach the destination



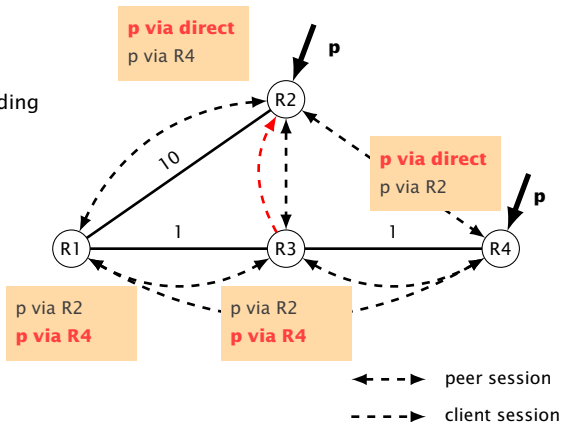
As for the IGP, reconfiguring BGP can create problems

Both R1 and R3 forwards traffic to R4



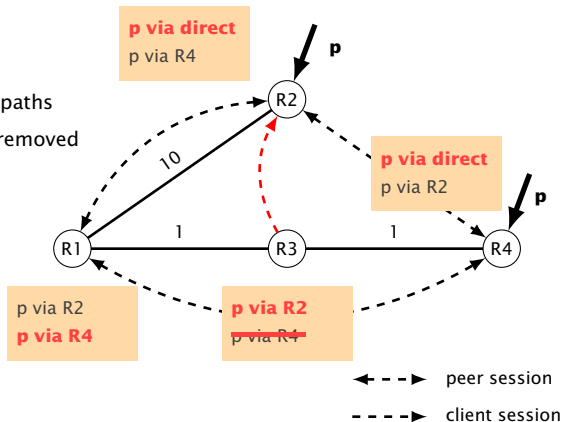
As for the IGP, reconfiguring BGP can create problems

R3 is reconfigured first by adding a client session towards R2



## As for the IGP, reconfiguring BGP can create problems

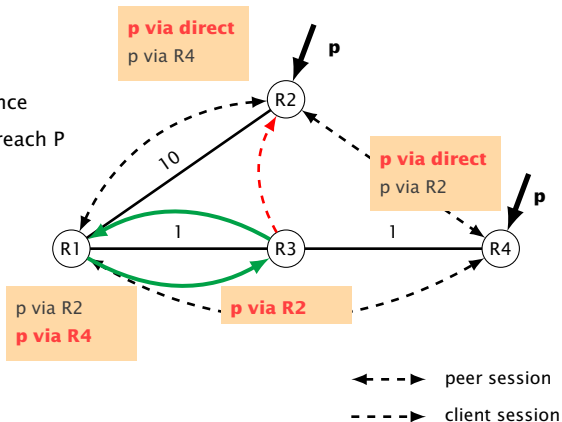
R3 loses the visibility of R4's paths  
when the initial sessions are removed



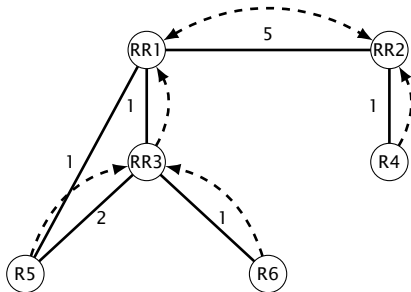


## As for the IGP, reconfiguring BGP can create problems

A forwarding loop appears since  
R3 and R4 use each other to reach P



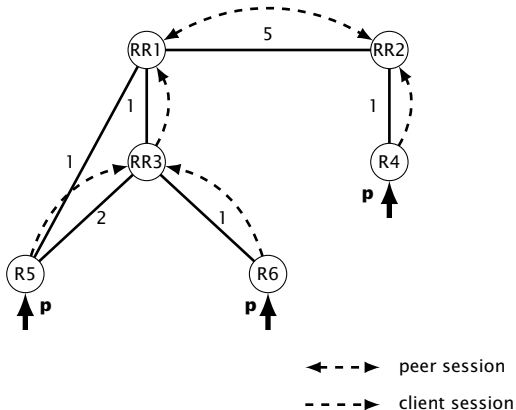
Adding iBGP sessions locally  
can worsen the visibility globally



← - - - → peer session  
- - - → client session

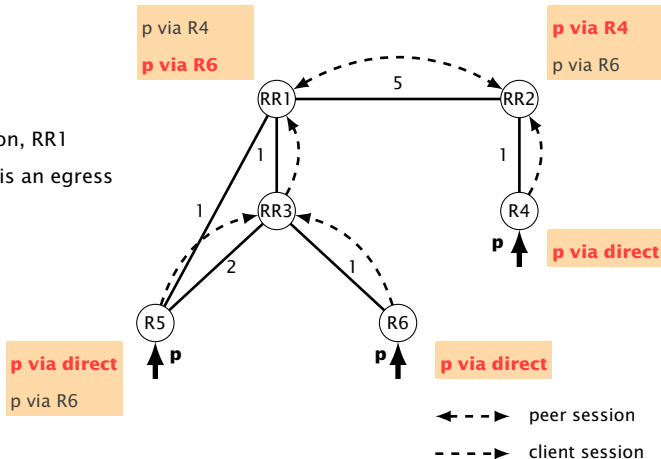
## Adding iBGP sessions locally can worsen the visibility globally

The same prefix is learned  
on R4, R5 and R6



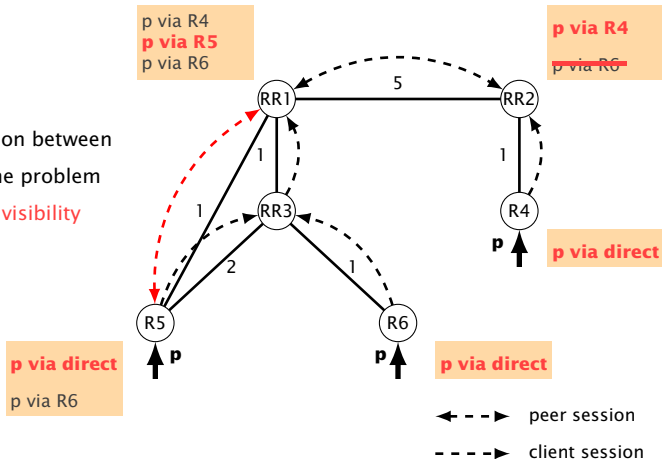
# Adding iBGP sessions locally can worsen the visibility globally

Due to route-reflection, RR1  
is not aware that R5 is an egress



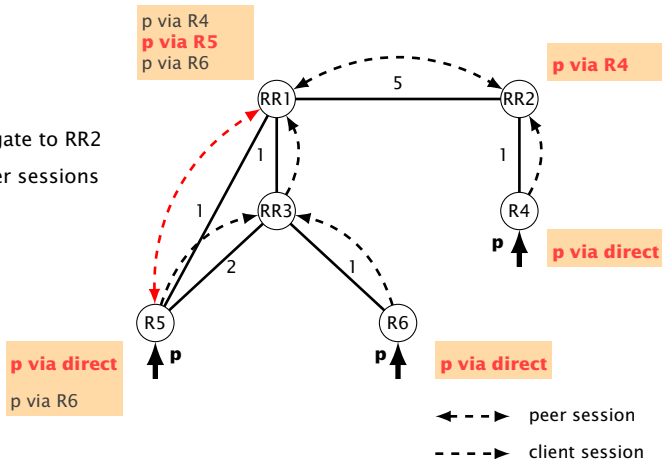
# Adding iBGP sessions locally can worsen the visibility globally

Adding an iBGP session between  
R5 and RR1 solves the problem  
... but worsen RR2's visibility



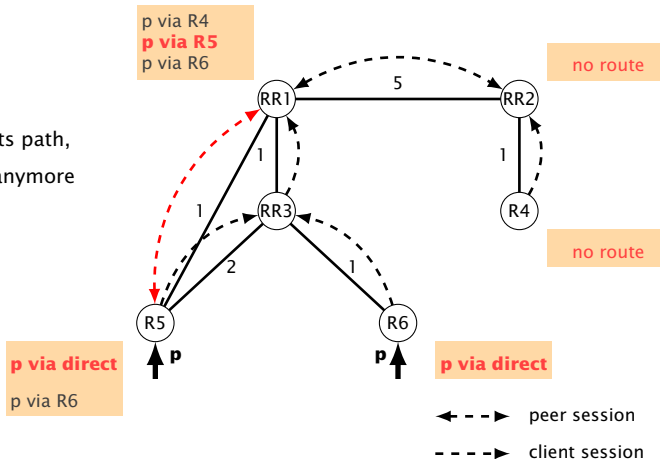
## Adding iBGP sessions locally can worsen the visibility globally

RR1 does not propagate to RR2  
paths learned on peer sessions



# Adding iBGP sessions locally can worsen the visibility globally

**Worse!** If R4 loses its path,  
RR2 cannot reach P anymore



# Reconfiguring BGP relying only on the protocol is quite challenging

Checking if adding (removing)  
an iBGP session will trigger

- forwarding loops
- routing instabilities
- visibility issues
- traffic shifts

is **computationally intractable**



# Leveraging other technologies helps in achieving seamless BGP reconfiguration

Use the BGP/MPLS VPNs machinery

- Deploy the different configurations in different VRFs
- Migrate by switching the VRF used to forward traffic

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A lot of research challenges still remain ...

# Seamless Network-Wide Migrations



IGP migrations

Ordering matters

BGP migrations

Ongoing work

3

How can we help ?

Manage complexity

# By using our techniques, we can help you achieve lossless migrations

Based on the initial and the final configurations, we can

IGP

Compute the operational ordering you should follow in order to not loose any packets

BGP

Evaluate the damages of the reconfiguration and work on ad-hoc solutions by modeling the network behavior

Please, come and talk to me if you are interested

# Seamless Network-Wide Migrations



IGP migrations

Ordering matters

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Ongoing work

How can we help ?

Manage complexity

# Don't fear network reconfiguration, **adapt** the network to its environment

Seamless reconfigurations are possible  
but they require careful planning

Our techniques provide theoretical guarantees  
that every reconfiguration step is safe

Add flexibility to your network management  
seamlessly move to the current best configuration

# Seamless Network-wide Migrations towards more agile networking



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